

MATH 281A FALL 2016 HOMEWORK 3

Due date: Friday December 2

**Rules:** Write as efficiently as possible: Include all relevant points and think carefully what to write and what not. Use common sense to determine what is the appropriate amount of details for a course at this level. Quote any result from the lecture you are referring to; do not reprove the result. Each problem indicates maximum allowed length; this is usually much more than needed. If you type, do not use font smaller than 10pt.

**I will not grade any text that exceeds the specified length.**

**1. (1 page)** A function  $g : {}^{<\omega}\omega \rightarrow {}^{<\omega}\omega$  is called monotonic iff the following two conditions are satisfied.

- (a)  $s \subseteq t \implies g(s) \subseteq g(t)$ , and
- (b) for every  $x \in \mathcal{N}$  the lengths  $\text{lh}(g(x \upharpoonright n))$  converge to  $\omega$  as  $n$  converges to  $\omega$ .

Given a monotonic function  $g : {}^{<\omega}\omega \rightarrow {}^{<\omega}\omega$  we can define a function  $g^* : \mathcal{N} \rightarrow \mathcal{N}$  by

$$g^*(x) = \bigcup_n g(x \upharpoonright n).$$

Prove the following.

- (i) If  $g : {}^{<\omega}\omega \rightarrow {}^{<\omega}\omega$  is monotonic then  $g^* : \mathcal{N} \rightarrow \mathcal{N}$  is continuous.
- (ii) If  $g : {}^{<\omega}\omega \rightarrow {}^{<\omega}\omega$  is monotonic then  $g^* : \mathcal{N} \rightarrow \mathcal{N}$  is  $\Sigma_1^0$ -recursive in  $a_g$  where  $a_g \in \mathcal{N}$  is a natural coding of  $g$ .
- (iii) If  $f : \mathcal{N} \rightarrow \mathcal{N}$  is continuous then there is a monotonic  $g : {}^{<\omega}\omega \rightarrow {}^{<\omega}\omega$  such that  $f = g^*$ .

**Remark.** Recall that if  $\Gamma$  is a pointclass then a function  $f : \mathcal{N} \rightarrow \mathcal{N}$  is  $\Gamma$ -recursive iff the relation  $P^f \subseteq \mathcal{N} \times \omega$  defined by

$$P^f(x, s) \iff \text{SEQ}(s) \wedge s \subseteq f(x)$$

is in  $\Gamma$ .

Thus, every continuous function  $f : \mathcal{N} \rightarrow \mathcal{N}$  is  $\Sigma_1^0$ -recursive in some  $a \in \mathcal{N}$ . Consider the following analogy: The function  $g$  in (iii) relates to  $f$  in an analogous way as the restriction of  $f$  to rationals, should  $f$  be considered as a function from reals to reals.

**2. (1/3 page)** Recall that if  $A \subseteq \mathcal{N}$  is a  $\Pi_1^1$ -set then there is a recursive tree  $T$  on  $\omega \times \omega$  such that for every  $x \in \mathcal{N}$ ,

$$A(x) \iff \langle x \rangle \in \text{WO}.$$

We proved in the lecture that  $A$  is  $\Delta_1^1$  iff the function  $x \mapsto \text{otp}(\langle x \rangle)$  is bounded below  $\omega_1^{\text{CK}}$ .

Prove that there is a  $\Pi_1^1$ -set (lightface!)  $A \subseteq \mathcal{N}$  such that the values  $\text{otp}(\langle x \rangle)$  are unbounded in  $\omega_1$  (this is the true  $\omega_1$ !) as  $x$  ranges over  $A$ .

**3.** A pointclass  $\Gamma$  has the reduction property iff for any sets  $A, B \in \Gamma$  such that  $A, B \subseteq \mathcal{N}$  there are sets  $A^*, B^* \in \Gamma$  such that

- (r1)  $A^* \subseteq A$  and  $B^* \subseteq B$ .
- (r2)  $A^* \cap B^* = \emptyset$  and  $A^* \cup B^* = A \cup B$ .

A pointclass  $\Gamma$  has the separation property iff for any sets  $A, B \in \Gamma$  such that  $A, B \subseteq \mathcal{N}$  there are sets  $A', B' \in \Gamma$  such that

- (s1)  $A \subseteq A'$  and  $B \subseteq B'$ .
- (s2)  $A' \cap B' = A \cap B$  and  $A' \cup B' = \mathcal{N}$ .

Prove the following

- (a) **(5 lines)** If  $A, B \subseteq \mathcal{N}$  are in  $\Gamma$  and  $A \cap B = \emptyset$  then the separation property for  $A, B$  is equivalent to the statement that there is a set  $D \in \Delta = \Gamma \cap \check{\Gamma}$  such that  $A \subseteq D$  and  $B \cap D = \emptyset$ . (This is the form of separation property we considered in Math 280.)
- (b) **(8 lines)**  $\Gamma$  has the reduction property iff  $\check{\Gamma}$  has the separation property.
- (c) **(1/3 page)** If  $\Gamma$  is normed then  $\Gamma$  has the reduction property. Conclude that  $\Pi_1^1$  has the reduction property and  $\Sigma_1^1$  has the separation property. Similarly for relativized and boldface pointclasses. This gives a lightface version of the separation theorem for analytic sets we proved in Math 280.

Recall that  $\Gamma$  is  $\omega$ -parametrized iff there is a universal set  $G \in \Gamma$  such that  $G \subseteq \omega \times \mathcal{N}$  and for every  $P \in \Gamma$  there is  $e \in \omega$  such that  $P(x) \iff G(e, x)$  for all  $x \in \mathcal{N}$ .

- (d) **(2/3 page)** Prove that if  $\Gamma$  is  $\omega$ -parametrized then  $\Gamma$  cannot have both the reduction property and the separation property. In fact,  $\Gamma$  cannot have both the reduction property and the weak form of separation property for disjoint sets discussed in (a) above.

It follows that if  $n$  is such that  $\Sigma_n^1/\Pi_n^1$  is normed then  $\Pi_n^1/\Sigma_n^1$  is not normed. In particular,  $\Sigma_1^1$  and  $\Pi_2^1$  are not normed (and therefore not scaled) pointclasses.

**Hint.** For (d), fix a universal  $\Gamma$ -set  $G \subseteq \omega \times \mathcal{N}$ . Let  $A, B \subseteq \mathcal{N}$  be defined by

$$A(x) \equiv G(x(0), x) \quad \text{and} \quad B(x) \equiv G(x(1), x).$$

Let  $A^*, B^*$  reduce  $A, B$  and  $A', B'$  separate  $A^*, B^*$ . Let  $i, j \in \omega$  be such that

$$A'(x) \iff G(i, x) \quad \text{and} \quad B'(x) \iff G(j, x),$$

and let  $a \in \mathcal{N}$  be such that  $a(0) = j$  and  $a(1) = i$ . Since  $A' \cup B' = \mathcal{N}$  and  $A' \cap B' = \emptyset$ , the sequence  $a$  is an element of exactly one of  $A'$  or  $B'$ . Use this to get a contradiction.