

# COVERING THE HYPERCUBE, THE UNCERTAINTY PRINCIPLE, AND AN INTERPOLATION FORMULA

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**ABSTRACT.** We show that the minimal number of skewed hyperplanes that cover the hypercube  $\{0, 1\}^n$  is at least  $\frac{n}{2} + 1$ , and there are infinitely many  $n$ 's when the hypercube can be covered with  $n - \log_2(n) + 1$  skewed hyperplanes. The minimal covering problems are closely related to the uncertainty principle on the hypercube, where we also obtain an interpolation formula for multilinear polynomials on  $\mathbb{R}^n$  of degree less than  $\lfloor n/m \rfloor$  by showing that its coefficients corresponding to the largest monomials can be represented as a linear combination of values of the polynomial over the points  $\{0, 1\}^n$  whose Hamming weights are divisible by  $m$ .

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## 1. INTRODUCTION

**1.1. Covering the hypercube.** How many affine hyperplanes are needed to cover the hypercube  $\{-1, 1\}^n$ ? Notice that two affine hyperplanes  $x_1 = -1$  and  $x_1 = 1$  cover the hypercube, and clearly this is the minimal number. However, if one requires that the affine hyperplanes are skewed, i.e.,  $a_1x_1 + \dots + a_nx_n + b = 0$  with all  $a_1, \dots, a_n \neq 0$ , then the problem becomes challenging<sup>1</sup>.

It follows from Littlewood–Offord inequalities that any skewed hyperplane covers at most  $n^{-1/2}$  fraction of the points in  $\{-1, 1\}^n$  (up to a universal constant factor), therefore, one needs at least  $\Omega(n^{1/2})$  skewed hyperplanes to cover the hypercube. In [7], this lower bound was improved to  $\Omega(n^{0.51})$ , and recently in [4] to  $\Omega(n^{2/3} \log(n)^{-4/3})$  by the second named author of the present paper.

The family of  $n + 1$  hyperplanes,  $x_1 + \dots + x_n = 2k - n$  for all  $k = 0, \dots, n$ , covers the hypercube. In fact, if  $n$  is even, one can cover with  $n$  skewed hyperplanes just by replacing the two hyperplanes corresponding to  $k = 0$  and  $k = n$  in the previous example with one hyperplane  $x_1 + \dots + x_{n/2} - x_{n/2-1} - \dots - x_n = 0$ . Moreover, it follows from [1] that for even  $n$ , the upper bound  $n$  on the minimal cover is also a lower bound if one restricts the covering to the family of “regular” hyperplanes, i.e., the ones  $\varepsilon_1x_1 + \dots + \varepsilon_nx_n + b = 0$  with  $\varepsilon_j = \pm 1$  for all  $j = 1, \dots, n$ .

Looking at the results for the case of “regular” hyperplane cover in [1], one may suspect that in analogy to Littlewood–Offord problem the sharp lower bound on the minimal skew hyperplane cover should be  $n$ . Surprisingly, one can cover the hypercube  $\{-1, 1\}^5$  with the following 4 skewed hyperplanes

$$\begin{aligned} x_1 + x_2 + x_3 + x_4 + 2x_5 &= 0, \\ x_1 + x_2 + x_3 - x_4 + 2x_5 &= 0, \\ x_1 + x_2 + x_3 + x_4 - 2x_5 &= 0, \\ x_1 + x_2 + x_3 - x_4 - 2x_5 &= 0. \end{aligned}$$

Also the hypercube  $\{-1, 1\}^6$  can be covered with 5 skewed hyperplanes (see Section 2.3). In fact, one can cover the hypercube  $\{-1, 1\}^n$  with  $n - \log_2(n) + 1$  skewed hyperplanes for infinitely many  $n$ 's.

<sup>1</sup>In what follows we will omit the word affine and we will be referring to such hyperplanes as skewed hyperplanes.

**Proposition 1.** For any integer  $m \geq 1$  the hypercube  $\{-1, 1\}^{2^m+m-1}$  can be covered with  $2^m$  skewed hyperplanes.

This proposition shows that the skew hyperplane covering problem is genuinely different from the original “regular” problem solved in [1].

**Question 2.** What is the minimal number of skewed hyperplane cover of the hypercube  $\{-1, 1\}^n$ ?

We prove the following lower bound.

**Theorem 3.** The minimal number of skewed hyperplane cover of  $\{-1, 1\}^n$  is at least  $\frac{n}{2} + 1$ .

There is a close relation between the minimal hyperplane covering problem and the uncertainty principle on the hypercube. Let  $p(x)$  be a polynomial on  $\mathbb{R}^n$ , and let  $\text{supp}(p) = \{x \in \mathbb{R}^n : p(x) \neq 0\}$ . Under what conditions on  $\text{supp}(p) \cap \{-1, 1\}^n$  and  $\deg(p)$  does it follow that  $p(x) \equiv 0$  on  $\{-1, 1\}^n$ ?

It turns out that the support of a nonzero low degree polynomial cannot be contained in a skewed hyperplane:

**Theorem 4** (Linial–Radhakrishnan [5]). If  $\deg(p) < \frac{n}{2}$  and  $\text{supp}(p) \cap \{-1, 1\}^n$  belongs to a skewed hyperplane, then  $p(x) \equiv 0$  on  $\{-1, 1\}^n$ .

Observe that Theorem 3 follows from Theorem 4. Indeed, let  $H_1, \dots, H_{k+1}$  be a minimal skew hyperplane cover of  $\{-1, 1\}^n$ . If  $H_j$ 's are given via equations  $\ell_j(x) = a_{1j}x_1 + \dots + a_{nj}x_n + b_j = 0$ , for all  $j = 1, \dots, k+1$ , then it follows that  $p(x)\ell_{k+1}(x) \equiv 0$  on  $\{-1, 1\}^n$ , where  $p(x) = \prod_{j=1}^k \ell_j(x)$  is a not identically zero polynomial on  $\{-1, 1\}^n$  of degree at most  $k$ . Hence,  $\text{supp}(p) \cap \{-1, 1\}^n$  belongs to  $H_{k+1}$  and Theorem 4 implies that  $k \geq n/2$ .

After the current paper was completed, independently and concurrently the paper [6] appeared on arXiv where Theorem 3 was derived from Theorem 4 proved in [5] (see Lemma 2 in [5]). The proof of Theorem 4 in [5] in turn is based on either Combinatorial Nullstellensatz or the spectral properties of the Johnson graph (the authors [5] attribute the nonsingularity of the Johnson graph to [3]). Our proof of Theorem 4, given in Section 2.1, is simple and self-contained.

**1.2. An interpolation formula.** In [1] sharp lower bound  $n$  on the minimal number of “regular” hyperplane cover of the  $n$  dimensional hypercube (for even  $n$ ) was based on the following technical observation: if a multilinear polynomial  $p(x)$  vanishes on all those points of  $\{-1, 1\}^n$  which have even number of 1’s in its coordinates, and  $\deg(p) < n/2$ , then  $p$  is identically zero (see Lemma 2.1 in [1]). This observation suggests that perhaps the coefficients of the multilinear polynomials of small degree can be reconstructed by its values at sparse points of  $\{-1, 1\}^n$ . The goal of this section is to obtain such an interpolation formula.

Recall that any function  $f : \{-1, 1\}^n \mapsto X$ , where  $X$  is a normed space, has Fourier–Walsh representation

$$f(x) = \sum_{S \subset \{1, \dots, n\}} \widehat{f}(S) x^S, \tag{1}$$

for some  $\widehat{f}(S) \in X$ , where  $x = (x_1, \dots, x_n)$ ,  $x^S = \prod_{j \in S} x_j$  and  $x^\emptyset = 1$ . We say that  $f$  has degree  $\deg(f)$  if  $\widehat{f}(S) = 0$  for all  $S \subset \{1, \dots, n\}$  with  $|S| > \deg(f)$ , and there exists a subset  $S$  of cardinality  $\deg(f)$  such that  $\widehat{f}(S) \neq 0$ .

**Definition 5.** For any integer  $m > 1$  the symbol  $W(m)$  denotes the subset of  $\{-1, 1\}^n$  consisting of all points  $x = (x_1, \dots, x_n) \in \{-1, 1\}^n$  such that  $\#\{j : x_j = -1\}$  is divisible by  $m$ .

In this section we obtain the following interpolation formula

**Theorem 6.** Let  $f : \{-1, 1\}^n \mapsto X$  and let  $m \geq 2$  be an integer divisible by 2 such that

$$\deg(f) \leq \frac{n}{m} - \frac{1}{2}. \quad (2)$$

Then for any  $S \subset [n]$  with  $|S| = \deg(f)$ , there exists a probability measure  $d\mu(x)$  supported on  $W(m)$  and a sign function  $h : W(m) \mapsto \{-1, 1\}$  such that

$$\widehat{f}(S) = \int_{W(m)} h(x)f(x)d\mu(x) \quad (3)$$

Both  $d\mu$  and  $h$  depend only on  $S, m, \deg(f), n$ .

The next corollary follows from the theorem

**Corollary 7.** If  $f : \{-1, 1\}^n \mapsto X$  vanishes on a set  $W(m)$  for some even integer  $m$  satisfying (2), then  $f \equiv 0$ .

**Remark 8.** When  $m = 2$ , Corollary 7 is the classical result [1, Lemma 2.1].

**Remark 9.** In the proof of Theorem 6 both the measure  $d\mu$  and  $h(x)$  are constructed explicitly.

Notice that since  $\widehat{f}(S) = \mathbb{E}f(x)x^S$  then  $\|\widehat{f}(S)\| \leq \max_{x \in \{-1, 1\}^n} \|f(x)\|$ . However, if  $f$  has low degree, then  $\max_{x \in \{-1, 1\}^n} \|f(x)\|$  can be replaced by a maximum over sparse family of points of  $\{-1, 1\}^n$  provided that  $|S| = \deg(f)$ . Indeed, Theorem 6 gives the following

**Corollary 10.** Let  $f : \{-1, 1\}^n \mapsto X$ , where  $X$  is a normed space, be a function whose degree satisfies (2), then

$$\|\widehat{f}(S)\| \leq \max_{x \in W(m)} \|f(x)\|$$

for all  $S \subset \{1, \dots, n\}$  with  $|S| = \deg(f)$ .

## 2. PROOFS

**2.1. The proof of Theorem 4.** Denote  $[n] := \{1, \dots, n\}$ . Every polynomial  $p(x)$  of degree  $d$ , when restricted to  $\{-1, 1\}^n$ , can be written as  $f(x) = \sum_{|S| \leq d} c_S x^S$  for some  $c_S \in \mathbb{R}$ . The assumption in Theorem 4 that the support of  $f$  is contained within a skewed hyperplane means that

$$\forall x \in \{-1, 1\}^n : (a_1 x_1 + \dots + a_n x_n + b) \sum_{|S| \leq d} c_S x^S = 0, \quad (4)$$

where

$$\forall i : a_i \neq 0. \quad (5)$$

When expanding (4), the right hand side means that all coefficients of the monomials  $x^T$  must vanish, in particular for degree  $d + 1$  monomials. This means that

$$\sum_{j \in T} a_j c_{T \setminus j} = 0 \quad (6)$$

for all  $T \subseteq [n]$  with  $|T| = d + 1$ . We can view (6) as a system of linear equations  $Ac = 0$  in  $c = (c_S : S \subset [n], |S| = d)$ . It suffices to prove

**Lemma 11.** We have  $\text{Ker}(A) = 0$  as long as  $n \geq 2d + 1$ .

The theorem follows from the lemma as follows: The lemma implies that in (4) we have  $c_S = 0$  for all  $S \subset [n]$  with  $|S| = d$ . This means that  $p(x)$  is in fact a polynomial of degree  $d - 1$ . Since  $d$  was merely defined as the degree of  $p$  (and assumed to satisfy  $2d + 1 \leq n$ ), we can repeat and deduce  $p$  is of degree  $d - 2$  and similarly of degree 0. Once  $n \geq 2$  a hyperplane can not cover the entire hypercube, so  $p$  must be the zero polynomial, concluding the proof of Theorem 4.

*Proof of Lemma 11.* Note that it is sufficient to prove the lemma only for  $n = 2d + 1$ , and it would follow for any  $n \geq 2d + 1$ . To see this, let  $n > 2d + 1$  and let  $S \subset [n]$  with  $|S| = d$ ; we must show  $c_S = 0$ . Fix a set  $N \subset [n]$  with  $S \subset N$  and  $|N| = 2d + 1$ , and focus on equations (6) for  $T \subset N$ . From the  $n = 2d + 1$  case we conclude that all involved variables  $c_{S'}$ , with  $S' \subset N$  are 0, and in particular  $c_S = 0$ .

Next, we prove the lemma by induction on  $d$  for the  $n = 2d + 1$  case.

**Base case:** When  $d = 1$  and  $n = 3$ , Equation (6) applied on the sets  $T = \{2, 3\}, \{1, 3\}, \{1, 2\}$  yields:

$$\begin{aligned} a_3 c_2 + a_2 c_3 &= 0 \\ a_1 c_3 + a_3 c_1 &= 0 \\ a_1 c_2 + a_2 c_1 &= 0, \end{aligned}$$

Or in matrix form:

$$\begin{bmatrix} 0 & a_3 & a_2 \\ a_3 & 0 & a_1 \\ a_2 & a_1 & 0 \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ c_3 \end{bmatrix} = 0.$$

The determinant of this  $3 \times 3$  matrix equals  $2a_1 a_2 a_3$ , which is nonzero by (5). We learn that  $c_i = 0$  and the lemma follows in this case.

**Inductive step:** Assume that the lemma holds for  $d - 1$ , we prove it for  $d$  and  $n = 2d + 1$ . Let  $c$  be a solution to (6). In order to complete the induction step, we must show that  $c = 0$ .

Note that the induction hypothesis applied for  $d - 1$  and  $n - 2$  (that has  $n - 2 = 2(d - 1) + 1$ ) implies that  $\bar{c} = 0$  is the unique solution to the system of equations

$$a_{i_1} \bar{c}_{i_2, i_3, \dots, i_d} + a_{i_2} \bar{c}_{i_1, i_3, \dots, i_d} + \dots + a_{i_d} \bar{c}_{i_1, i_2, \dots, i_{d-1}} = 0 \quad \forall \{i_1, \dots, i_d\} \subset [n - 2], \quad (7)$$

where  $\{i_1, \dots, i_d\}$  ranges over all subsets of size  $d$  of  $[n - 2]$ .

Fix  $j := i_{d+1} \in \{n - 1, n\}$ , and consider all those linear equations (6) in  $c$  that arise for  $T = \{i_1, \dots, i_d, j\}$  where  $\{i_1, \dots, i_d\} \subset [n - 2]$ :

$$a_{i_1} c_{i_2, i_3, \dots, i_d, j} + a_{i_2} c_{i_1, i_3, \dots, i_d, j} + \dots + a_{i_d} c_{i_1, i_2, \dots, i_{d-1}, j} + a_j c_{i_1, i_2, \dots, i_d} = 0 \quad \forall \{i_1, \dots, i_d\} \subset [n - 2].$$

Moving the last term into the right hand side, we get:

$$a_{i_1} c_{i_2, i_3, \dots, i_d, j} + a_{i_2} c_{i_1, i_3, \dots, i_d, j} + \dots + a_{i_d} c_{i_1, i_2, \dots, i_{d-1}, j} = -a_j c_{i_1, i_2, \dots, i_d} \quad \forall \{i_1, \dots, i_d\} \subset [n - 2]. \quad (8)$$

Compare this system with the one in (7). This is the same system, up to a relabeling of the variables, and different right hand side. Let  $A_{d-1}$  be the matrix defining the  $d - 1$  case, then (8) can be written as

$$A_{d-1} \bar{c} = -a_j u,$$

where  $u$  is the vector of variables  $c_{i_1, i_2, \dots, i_d}$  and  $\bar{c}$  is the vector of variables  $c_{i_1, i_2, \dots, i_d, j}$  for  $\{i_1, i_2, \dots, i_d\} \subset [n - 2]$ . Note that since  $\binom{2d-1}{d-1} = \binom{2d-1}{d}$ , the matrix  $A_{d-1}$  is square, and is hence invertible by the induction hypothesis. In particular,

$$\bar{c} = -a_j A_{d-1}^{-1} u.$$

Note that on the right hand side, the only thing that depends on  $j$  is  $a_j$ . That is, both  $A_{d-1}$  and  $u$  do not depend on whether  $j = n - 1$  or  $j = n$ . By comparing the two options  $j = n - 1$ , and  $j' = n$ , we conclude that for all  $i_1, i_2, \dots, i_d \subset [n - 2]$ ,

$$c_{i_1, i_2, \dots, i_d, j} / a_j = c_{i_1, i_2, \dots, i_d, j'} / a_{j'}. \quad (9)$$

Note that in (9), we can choose the indices  $i_1, \dots, i_d, j', j$  arbitrarily so long as they are distinct.

**Claim 12.** Equation (9) implies the existence of a single constant  $K \in \mathbb{R}$  such that

$$c_S = K \prod_{i \in S} a_i, \quad (10)$$

for all  $S \subset [n]$  with  $|S| = d$ .

Plugging the formula for  $c_S$  from Claim 12 into the system of equations (6), we get that for any  $T \subset [n]$  with  $|T| = d + 1$ ,

$$0 = \sum_j a_j c_{T \setminus \{j\}} = \sum_{j \in T} a_j \left( K \prod_{i \in T \setminus \{j\}} a_i \right) = \sum_{j \in T} K \prod_{i \in T} a_i = (d + 1) K \prod_{i \in T} a_i$$

In (5), we assumed  $a_i \neq 0$  for all  $i$ , so it follows that  $K = 0$ . and in particular  $c_S = 0$  for all  $S \subset [n]$ . The induction step and the lemma follows.

*Proof of Claim 12.* Let  $I \subset [n]$  and  $J, K \subset [n] \setminus I$  be sets with  $|J| = |K| = d - |I|$ . We prove by induction on the size of  $J$  and  $K$  that

$$c_{I \cup J} / \prod_{j \in J} a_j = c_{I \cup K} / \prod_{k \in K} a_k. \quad (11)$$

When  $J, K$  are of size  $d$ , and  $I = \emptyset$ , then we derive (10) and the proof is complete.

**Base case:** When  $J, K$  are of size 0, (11) is obvious as  $J = K$ .

**Inductive step:** Let  $J, K \subset [n]$  be of size  $\ell \geq 1$  and let  $I \subset [n] \setminus (J \cup K)$  be of size  $d - \ell$ . We prove (11). Take  $j \in J$  and  $k \in K$  and denote  $I' = I \cup \{j\}$ ,  $J' = J \setminus \{j\}$  and  $K' = K \setminus \{k\}$ . Then (11) follows:

$$\begin{aligned} c_{I \cup J} / \prod_{i \in J} a_i &\stackrel{I \cup J = I' \cup J'}{=} \frac{1}{a_j} c_{I' \cup J'} / \prod_{i \in J'} a_i \\ &\stackrel{\text{induction}}{=} \frac{1}{a_j} c_{I' \cup K'} / \prod_{i \in K'} a_i \\ &\stackrel{(9)}{=} \frac{1}{a_k} c_{I \cup K} / \prod_{i \in K'} a_i \\ &= c_{I \cup K} / \prod_{i \in K} a_i. \end{aligned}$$

□

□

**2.2. The proof of Theorem 6.** Let  $d = \deg(f)$  and assume that  $S = \{m, 2m, \dots, dm\}$  (any  $S$  with  $|S| = d$  is obtained by relabeling of variables). Recall Equation (1) defining the Fourier representation: for any  $u = (u_1, \dots, u_n) \in \{-1, 1\}^n$  we have

$$f(u) = \sum_{|S| \leq d} \widehat{f}(S) u^S \quad \text{where} \quad u^S = \prod_{j \in S} u_j.$$

For  $(y_1, \dots, y_d) \in \{-1, 1\}^d$ , and  $(x_1, \dots, x_n) \in \{-1, 1\}^n$ , we define  $y \circ x \in \{-1, 1\}^n$  by splitting the vector  $x$  into disjoint sets of indices  $I_1, I_2, \dots, I_d, I_{\text{extra}}, I_{\text{rest}}$ , where

$$\begin{aligned} I_1 &= (1, \dots, m), \\ I_2 &= (m + 1, \dots, 2m), \\ &\dots \\ I_d &= ((d - 1)m + 1, \dots, dm), \\ I_{\text{extra}} &= (md + 1, \dots, md + m/2), \\ I_{\text{rest}} &= (md + m/2 + 1, \dots, n). \end{aligned}$$

Then, we define

$$y_j x_{I_j} := (y_j x_{(j-1)m+1}, y_j x_{(j-1)m+2}, \dots, y_j x_{jm}),$$

and finally,

$$\begin{aligned}
y \circ x &:= (y_1 x_{I_1}, \dots, y_d x_{I_d}, x_{I_{\text{extra}}}, x_{I_{\text{rest}}}) \\
&= (y_1 x_1, y_1 x_2, \dots, y_1 x_m, \\
&\quad y_2 x_{m+1}, y_2 x_{m+2}, \dots, y_2 x_{2m}, \\
&\quad \vdots \\
&\quad y_d x_{(d-1)m+1}, y_d x_{(d-1)m+2}, \dots, y_d x_{md}, \\
&\quad x_{md+1}, x_{md+2}, \dots, x_{md+m-\frac{m}{2}}, \dots, x_n)
\end{aligned}$$

Note that the variables in  $I_{\text{extra}}$  and  $I_{\text{rest}}$  are unchanged. The coordinates  $I_{\text{rest}}$  do not play a role in the proof (and may be empty, e.g. if we have equality in (2)), but the coordinates  $I_{\text{extra}}$  have the important role of ‘‘parity’’ in the proof.

**Claim 13.** *There exists a distribution  $\mathcal{D}$  for  $x$  (on  $\{-1, 1\}^n$ ) such that*

$$\widehat{f}(S) = \mathbb{E}_{x \sim \mathcal{D}} \mathbb{E}_{y \sim \text{unif}(\{-1, 1\}^d)} [f(y \circ x) \cdot y_1 \cdots y_d], \quad (12)$$

and moreover,

$$y \circ x \in W(m) \quad (13)$$

for all  $y \in \{-1, 1\}^d$  and all  $x \in \text{supp}(\mathcal{D})$ .  $\mathcal{D}$  depends only on  $d, m, n$  but not on  $f$ .

Claim 13 concludes the proof of (3) by using the sign function  $h = y_1 \cdots y_d$  and the measure  $d\mu$  depicting the distribution<sup>2</sup> of  $y \circ x$  where  $x \sim D$  and  $y \sim \text{unif}(\{-1, 1\}^d)$ .

**Proof of Claim 13.** Observe the formula

$$\mathbb{E}_{y \sim \text{unif}(\{-1, 1\}^d)} [f(y \circ x) y_1 \cdots y_d] = \sum_{T: \forall j: |T \cap I_j|=1} \widehat{f}(T) x^T. \quad (14)$$

To verify (14), we expand  $f(y \circ x)$  on the left hand side as  $\sum_T \widehat{f}(T)(y \circ x)^T$ . The number of times  $y_j$  appears in that expression is exactly  $|T \cap I_j|$ . If  $T \cap I_j$  is empty for some  $j$ , then  $(y \circ x)^T$  does not depend on  $y_j$ , and the multiplication by  $y_1 \cdots y_d$  in (14) zeroes out the term  $\widehat{f}(T)(y \circ x)^T$ . Hence relevant terms are only those with  $|T \cap I_j| \geq 1$  for all  $j = 1, \dots, d$ . But since  $f$  is of degree  $d$  to begin with, we must have  $|T \cap I_j| = 1$  for all  $j = 1, \dots, d$ .

**The distribution  $\mathcal{D}$ .** We describe how each chunk  $x_{I_j}$  of  $x \sim \mathcal{D}$  is drawn. All chunks are drawn independently of the other chunks, except for  $x_{I_{\text{extra}}}$  which is chosen last.

- $x_{I_j}$  for  $j = 1, \dots, d$ :

$$\Pr[x_{I_j} = (z_1, \dots, z_m)] = \begin{cases} 1/m & \text{if } z_1 = \cdots = z_m = 1, \\ \frac{1}{2^{\binom{m-2}{m/2-1}}} & \text{if } z_m = 1 \text{ and exactly } m/2 \text{ among } z_1, \dots, z_{m-1} \text{ are equal } -1, \\ 0 & \text{otherwise.} \end{cases}$$

Note that the sum over all probabilities is 1.

- $x_{I_{\text{rest}}}$ : we set  $x_{I_{\text{rest}}} = (1, 1, \dots, 1)$  always.
- $x_{I_{\text{extra}}}$ : Count the total number  $s$  of  $-1$ 's in all chunks  $x_{I_1}, x_{I_2}, \dots, x_{I_d}$ . Define

$$x_{I_{\text{extra}}} = \begin{cases} (1, \dots, 1) & \text{if } m|s, \\ (-1, \dots, -1) & \text{otherwise.} \end{cases} \quad (15)$$

<sup>2</sup>Note that here we define  $h$  as a function of  $y$  while the measure  $d\mu$  is of  $\{-1, 1\}^n$ . We claim that (12) implies that  $y_1 \cdots y_d$  is uniquely determined from  $y \circ x$  for all  $y \circ x$  having positive probability. To see this, note that formula (12) does not depend on  $f$ , yet for  $\widetilde{f}(z) = z^S$  it has 1 on the LHS while the RHS is bounded by 1 by the triangle inequality. This means  $\widetilde{f}(y \circ x) = y_1 \cdots y_d$ . Consequently, (3) holds with  $h = y_1 \cdots y_d = \widetilde{f}(y \circ x)$ , which is a function of  $y \circ x$ .

Observe that necessarily  $s$  is divisible by  $m/2$ , since each choice of  $x_{I_j}$  adds either 0 or  $m/2$  to  $s$ . For this reason, (15) defines  $x$  in such a way that  $x \in W(m)$ . Furthermore, for all  $y \in \{-1, 1\}^d$  we have  $y \circ x \in W(m)$  essentially because signs do not matter modulo 2.

Finally, in order to deduce (12) from (14), we must check that for all  $T \subseteq \{1, \dots, n\}$  with  $\forall j: |T \cap I_j| = 1$  we have

$$\mathbb{E}_{x \sim \mathcal{D}} [x^T] = \begin{cases} 1 & \text{if } T = S, \\ 0 & \text{otherwise.} \end{cases} \quad (16)$$

The case  $T = S$  is immediate, since by design  $x_{jm} = 1$  for all  $j = 1, \dots, d$ .

Suppose  $T \neq S$ . Focus on a particular coordinate  $t \in T$  with  $m \nmid t$  and let  $j \in \{1, \dots, d\}$  be the index with  $t \in I_j$ . Since  $T \cap I_j = \{t\}$ , we have that  $x^T$  is  $x_t \cdot x^{T \setminus \{t\}}$ . When  $x \sim \mathcal{D}$ ,  $x^{T \setminus \{t\}}$  is a random variable independent of  $x_t$ , as we draw the different chunks independently. Hence  $\mathbb{E}_{x \sim \mathcal{D}} [x^T] = \mathbb{E}_{x \sim \mathcal{D}} [x_t] \cdot \mathbb{E}_{x \sim \mathcal{D}} [x^{T \setminus \{t\}}]$ . In order to deduce (16) and finish the proof we just need to check that  $\mathbb{E}_{x \sim \mathcal{D}} [x_t] = 0$ . Indeed, by definition, the probability that  $x_t = 1$  is  $1/m + \frac{\binom{m-2}{m/2}}{2 \binom{m-2}{m/2-1}}$ , that is, either  $x_{I_j}$  is all 1's, or we need to choose  $m/2$  locations for  $-1$ 's in  $I_j$  out of  $I_j \setminus \{t, mj\}$ . This probability is  $1/2$ , validating  $\mathbb{E}_{x \sim \mathcal{D}} [x_t] = 0$ , concluding the proof.

**2.3. The proof of Proposition 1: how to cover the hypercube efficiently.** Consider  $2^m$  skewed hyperplanes

$$\sum_{j=1}^{2^m-1} x_j + \sum_{j=0}^{m-1} \pm 2^j x_{2^{m+j}} = 0.$$

Since any odd integer  $k$ ,  $-(2^m - 1) \leq k \leq 2^m - 1$  can be written as a sum  $\sum_{j=0}^{m-1} \pm 2^j$  for some choice of signs  $\pm$ , it follows that these hyperplanes cover the cube  $\{-1, 1\}^n$  with  $n = 2^m + m - 1$ .  $\square$

There are other examples that are not produced by the construction above. In particular, for  $n = 6$  the union of the following 5 skewed hyperplanes

$$\begin{aligned} x_1 - x_2 + 2x_3 + x_4 + x_5 + 2x_6 &= 0, \\ x_1 - x_2 + x_3 + x_4 + x_5 - x_6 &= 0, \\ x_1 - x_2 - x_3 + 2x_4 - 2x_5 + x_6 &= 0, \\ x_1 + x_2 + x_3 + x_4 + x_5 - x_6 &= 0, \\ x_1 - x_2 - 3x_3 + x_4 + x_5 - x_6 &= 0. \end{aligned}$$

cover the hypercube  $\{-1, 1\}^6$ .

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