## Non-Asymptotic Theory of Random Matrices Lecture 4: Dimension Reduction

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### 1 Introduction

Consider the set  $X = \{n \text{ points in } \mathbb{R}^N\}$  where both n and N are large. Our goal will be to reduce N so as to represent X in a space of dimension  $k \ll N$ . One motivation for this comes from computer science where X is a data structure.

#### Observation 1.

- 1. We can always take N=n, since  $\dim(\operatorname{span}(X))=n$  and  $X\subset \operatorname{span}(X)$ .
- 2. For N < n we may lose linear independence.

Goal: We want to construct a map  $T: \mathbb{R}^N \to \mathbb{R}^k$ , k << N such that pairwise distances in X are approximately preserved. That is, for all  $x, y \in X$  there exists  $\varepsilon > 0$  such that

$$(1-\varepsilon)\|x-y\|_2 < \|Tx-Ty\|_2 < (1+\varepsilon)\|x-y\|_2$$

Such a map is called an  $\varepsilon$ -embedding of our data structure X into  $\mathbb{R}^k$  (more precisely into  $\ell_2^k$ ). For us, T will be the linear map associated with a random matrix.

**Theorem 2** (Johnson-Lindenstrauss (J-L) Flattening Lemma [3]). Let X be an n-point set in a Hilbert space, and suppose  $\varepsilon \in (0,1)$ . Then there exists an  $\varepsilon$ -embedding of X into  $\ell_2^k$ , for some  $k \leq C\varepsilon^{-2}\log n$  and some constant C > 0

**Example 3.** To represent n points in  $\mathbb{R}^n$  in a computer we need to store  $n^2$  numbers. However, by the J-L Flattening Lemma we can store only  $O(n \log n)$  numbers and still reconstruct all distances within  $\varepsilon$ -error.

All known embeddings satisfying Theorem 2 are given by random matrices T such as

- Gaussian
- Bernoulli (Achlioptas [1])
- Orthogonal projections (J-L)

### 1.1 Random Projections in $\mathbb{R}^n$

- Random rotations (= random orthogonal matrices)
- Orthogonal group  $O_n = \{ \text{rotations in } \mathbb{R}^n \} = \{ \text{orthogonal } n \times n \text{ matrices} \}$

Here,  $O_n$  is a probability space with probability measure called "Haar Measure" (from topological group theory).

#### 1.2 Haar Measure

Suppose that M is a compact metric space (such as a sphere in  $\mathbb{R}^n$ ), and that G is a group of isometries of M.

Theorem 4 (Haar Measure).

- 1. There exists a Borel probability measure  $\mu$  on M which is invariant under G. That is,  $\mu(S) = \mu(gS)$  for all  $g \in G$  and  $S \subset M$ .
- 2. If G is transitive, then the Haar measure is unique. Here, "transitive" means for all  $x, y \in X$ , there exists  $g \in G$  such that gx = y.

*Proof:* See §1 of Milman-Schechtmann [7] for a simple 2-page proof.

**Example 5.** Let  $M = S^{n-1}$ ,  $G = O_n$ . Then  $\mu = usual$  Lebesgue measure.

**Example 6.** Let  $M = G = O_n$ . Here the metric on a rotation comes from the Hilbert-Schmidt norm. In this sense  $n \times n$  matrices can be viewed as "vectors" in  $\mathbb{R}^{n^2}$ , and the Hilbert-Schmidt norm provides a Euclidean distance between these matrices. This gives us a Haar measure on  $O_n$ .

### 1.3 Random rotations, random orthogonal matrices

**Fact 7** (How to compute Haar Measure). For a given  $z \in S^{n-1}$  and for a given random rotation (with respect to Haar measure)  $U \in O_n$ , Uz is a random vector which is uniformly distributed on  $S^{n-1}$ .

*Proof.* We want to show for all  $A \subset S^{n-1}$  that we have  $\mathbb{P}(Uz \in A) = \sigma(A)$ , where  $\mathbb{P}$  denotes probability (with respect to Haar measure) and  $\sigma$  is the usual uniform measure. Note that both sides of this equation define a probability measure on  $S^{n-1}$  which are rotationally invariant. The uniqueness part of the Haar Measure theorem implies that they coincide.

#### Definition 8.

- Random k-subspace of  $\mathbb{R}^n$  = random rotation of  $\mathbb{R}^k \subset \mathbb{R}^n$ , denoted by  $U^*(\mathbb{R}^k)$
- Random Projection = orthogonal projection onto a random k-subspace, denoted by  $U^*P_kU$ , where  $P_k: \mathbb{R}^n \to \mathbb{R}^k$  is a projection (i.e.,  $P_kU =$ the first k rows of matrix U)

**Proof of J-L Lemma** (for  $\varepsilon$ -embedding given by random projections) Layout of proof: First we will show that J-L holds for one particular fixed pair  $x, y \in X$ . Then we will show that it holds collectively for all possible pairs in X.

**Step 1.** Fix an arbitrary pair  $x, y \in X$ , and let z = x - y.

**Lemma 9** ("Norm" of a random projection). Let  $z \in \mathbb{R}^n$  be fixed, and let P be a random projection in  $\mathbb{R}^n$  onto a k-subspace. Then

- 1.  $(\mathbb{E}||Pz||_2^2)^{1/2} = \sqrt{k/n}$ , where  $\mathbb{E}$  denotes expectation
- 2. For  $\varepsilon > 0$  we have that

$$(1-\varepsilon)\sqrt{k/n} \le ||Pz||_2 \le (1+\varepsilon)\sqrt{k/n}$$

holds with probability  $1 - 2e^{-k\varepsilon^2/2}$ 

*Proof.* 1. Note that the statement "Projecting a fixed z onto a random subspace" is equivalent to the statement "Projecting a random z onto a fixed subspace." Then with random rotation U we have

$$||Pz||_2 = ||U^*P_kUz||_2 = ||P_kUz||_2 = ||P_kx||_2$$

where x is a random vector distributed uniformly on  $S^{n-1}$ . Then

$$\mathbb{E} \|P_k x\|_2^2 = \mathbb{E} \sum_{j=1}^k x_j^2 = \sum_{j=1}^k \mathbb{E} x_j^2 = k \mathbb{E} x_1^2 = k/n.$$

2. For the next part of Lemma 9 we want to show

Failure Probability 
$$\leq \mathbb{P}\left(\left|\|Pz\|_2 - \sqrt{k/n}\right| > \varepsilon\sqrt{k/n}\right)$$
  
=  $\sigma\left(\left\{x \in S^{n-1} : \left|\|P_kx\|_2 - \sqrt{k/n}\right| > \varepsilon\sqrt{k/n}\right\}\right)$ .

Let  $f(x) = ||P_k x||_2$ , and recognize that  $f: S^{n-1} \to \mathbb{R}^n$  is a 1-Lipschitz function. From Part 1 of this lemma we have that

$$\left(\mathbb{E}f^2\right)^{1/2} = \sqrt{k/n}.$$

Let  $p = \sigma(\lbrace x \in S^{n-1} : |f(x) - (\mathbb{E}f^2)^{1/2}| > \varepsilon \sqrt{k/n} \rbrace)$ . Then

$$p \le 2e^{-n(\varepsilon\sqrt{k/n})^2/2} = 2e^{-k\varepsilon^2/2}$$

where we used the results of Concentration of Measure discussed in  $\S 2$  of Lecture 3.

We next normalize our embedding as  $T := \sqrt{n/k} P$ . By Lemma 9, we have for any fixed  $z \in X$  that

$$(1 - \varepsilon) \|z\|_2 \le \|Tz\|_2 \le (1 + \varepsilon) \|z\|_2$$
 (1)

holds with high probability. Thus for any fixed  $x, y \in X$  and putting z = x - y we have that the inequality (1) holds with probability  $1 - 2e^{-k\varepsilon^2/2}$ .

Now we consider the *collective ensemble*. Notice that there are no more than  $n^2$  pairs of  $x, y \in X$ . Then taking the union over all pairs we have the probability that (1) fails for some pair is less than or equal to

$$n^2 2e^{-k\varepsilon^2\!/2}$$

which is less than 1 if  $k \geq C\varepsilon^{-2}\log n$ . Hence with positive probability (1) holds with for all pairs. (End of proof for the J-L Flattening Lemma.)

**Example 10** (Argument that  $k \geq C(\varepsilon) \log n$  using volume point of view). Suppose for all  $x, y \in X$  that

$$1/2 \le ||x - y||_2 \le 1.$$

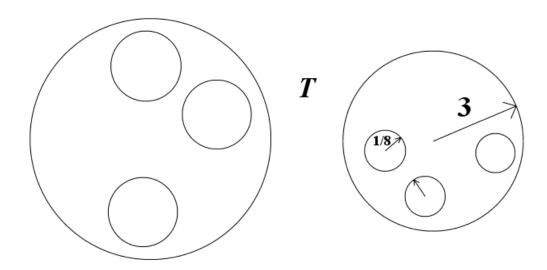
Now suppose that  $\varepsilon=1/10$ . Then by the J-L Flattening Lemma we can claim that

$$1/4 \le (1-\varepsilon)/2 \le ||Tx - Ty||_2 \le (1+\varepsilon) \le 2.$$

This tell us the following:

- The image of X under T is contained in a ball of radius 2 (i.e.,  $T(X) \subset 2B_2^k$ )
- There are n disjoint  $\frac{1}{8}$ -balls centered at points in  $T(X) \subset (2+\frac{1}{8})B_2^k \subset 3B_2^k$  if

$$\begin{array}{rcl} n \cdot \operatorname{Vol}(\frac{1}{8}B_2^k) & \leq & \operatorname{Vol}(3B_2^k) \\ n \cdot (\frac{1}{8})^k \operatorname{Vol}(B_2^k) & \leq & 3^k \operatorname{Vol}(B_2^k) \\ n & \leq & 24^k \\ \log n & \lesssim & k. \end{array}$$



The last example showed that the  $\log n$  factor in the J-L Flattening Lemma is sharp. The sharpness of the  $\varepsilon^{-2}$  factor is due to Alon and can be found in [6].

## 2 Related Work

Is there a version for

- $\ell_1$ ? No: Lee and Naor [4], Brinkman and Charikar [2].
- $\ell_{\infty}$ ? No.
- $\ell_p$  (for  $p \neq 1, 2, \infty$ )? Still open.

# 3 Related Readings

- Matousek [6] and [5]
- Vempala [8]

## References

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