

Non-Asymptotic Theory of Random Matrices

Lecture 6: Norm of a Random Matrix

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1 Introduction

Let A be an $m \times n$ random subgaussian matrix. That is, its entries x_{ij} are i.i.d. mean zero subgaussian random variables, or equivalently, they are independent copies of a mean zero subgaussian random variable X .

We know that the operator norm $\|A\| = s_1(A) = \max_{x \in S^{n-1}} \|Ax\|_2$, and we want to bound this with high probability.

- LOWER ESTIMATES (trivial): $\|A\| \geq \max(\|columns\|_2, \|rows\|_2)$

In particular, if A is a Bernoulli matrix, then $\|A\| \geq \max(\sqrt{n}, \sqrt{m})$.

- MATCHING UPPER ESTIMATES: norm of a random matrix \approx norms of rows or columns

– *Asymptotic*: Consider the mode when $n \rightarrow \infty$, $\frac{m}{n} \rightarrow y$. Let X be a random variable with $\mathbb{E}X^2 = 1$. (X need not be subgaussian in this case.) The best asymptotic result is that

$$\begin{aligned} \text{if } \mathbb{E}X^4 < \infty, \text{ then } \frac{\|A\|}{\sqrt{n}} &\rightarrow 1 + \sqrt{y} \quad a.s.; \\ \text{if } \mathbb{E}X^4 = \infty, \text{ then } \frac{\|A\|}{\sqrt{n}} &\rightarrow \infty \quad a.s.. \end{aligned}$$

For more details, see [5] [2] [1] [8].

– *Non-asymptotic*: We have the following main theorem in this lecture.

Theorem 1 (Upper Bound). *Let A be an $m \times n$ subgaussian random matrix. Then $\|A\| \leq C(\sqrt{n} + \sqrt{m})$ with probability $1 - \exp\{-c(m+n)\}$.*

Before we get to its proof, let us first pick up some useful tools.

2 Discretization of the Sphere

In the investigation of $\|A\|$, the difficulty is that the maximum “ $\max_{x \in S^{n-1}}$ ” is taken over infinitely many random variables. This is also called a random process. There are various tools to bound such random processes, from the ϵ -net method (simplest) to the majorizing measure method (hardest). We will work through the ϵ -net method in this lecture. For the majorizing measure method, see [3].

Idea: In order to reduce to a finite situation, we want first to discretize S^{n-1} , or in other words, to replace the sphere with a finite subset. This is possible because the sphere S^{n-1} is compact, and every compact set has a finite ϵ -net.

Definition 2. An ϵ -net of a subset K of a Banach space is a set \mathcal{N} such that $\forall z \in K, \exists x \in \mathcal{N} : \|z - x\| \leq \epsilon$.

Note: If \mathcal{N} is an ϵ -net, then every point of K is within distance ϵ from \mathcal{N} , and K is covered by balls of radius ϵ centered at $x \in \mathcal{N}$ (these balls are simply translates of $D := \epsilon B_X$).

Definition 3. The minimal cardinality of an ϵ -net of K is called the **covering number** of K by ϵ -balls, denoted by $N(K, D)$. In other words,

$$N(K, D) := \text{minimal number of translates of } D \text{ to cover } K.$$

Remark: The covering number gives a quantitative notion of compactness.

However, the covering number is hard to compute precisely even in the simplest case where we try to cover the unit disk by smaller disks in \mathbb{R}^2 . So we may want to make estimates. The good news is that relatively sharp estimates exist.

- LOWER ESTIMATE: $\text{Vol}(K) \leq N(K, D) \cdot \text{Vol}(D) \Rightarrow N(K, D) \geq \frac{\text{Vol}(K)}{\text{Vol}(D)}$.
- MATCHING UPPER ESTIMATE:

Proposition 4 (Covering Number). $N(K, D) \leq \frac{\text{Vol}(K) + \frac{1}{2}\text{Vol}(D)}{\text{Vol}(\frac{1}{2}D)}$

Proof. (constructive proof — greedy algorithm to find an ϵ -net of K)

We want to locate the centers in order to find the ϵ -net. Start with an arbitrary point $x_1 \in K$.

Then choose $x_2 \in K$: $\|x_2 - x_1\| > \epsilon$.

Then choose $x_3 \in K$: $\|x_3 - x_k\| > \epsilon, k = 1, 2$.

...

Then choose $x_N \in K$: $\|x_N - x_k\| > \epsilon, k = 1, \dots, N - 1$.

Stop if no more points are left.

Claim: $\mathcal{N} = \{x_1, \dots, x_N\}$ is an ϵ -net.

Suppose not. Then $\exists z \in K$ such that $\|z - x_k\| > \epsilon \quad \forall k$.

This contradicts the stopping criterion of the algorithm.

Then, if we shrink the radius of these balls to $\epsilon/2$, they will be disjoint (recall that each pair (x_i, x_j) is at least ϵ -apart). Thus, $\epsilon/2$ -balls (i.e. $\frac{1}{2}D$) centered at points of \mathcal{N} are disjoint and contained in $K + \frac{1}{2}D$. Therefore,

$$\text{Vol}(K + \frac{1}{2}D) \geq N \cdot \text{Vol}(\frac{1}{2}D),$$

and thus

$$N \leq \frac{\text{Vol}(K + \frac{1}{2}D)}{\text{Vol}(\frac{1}{2}D)}.$$

□

Remark: An upper bound for $\text{vol}(K + \frac{1}{2}D)$ was found by Bourgain-Milman [7].

Example 5. Let K be the unit ball B_X in \mathbb{R}^n . We want to cover the unit ball with smaller balls $D = \epsilon B_X$. Note that

$$K + \frac{1}{2}D = B_X + \frac{1}{2}\epsilon B_X = (1 + \frac{1}{2}\epsilon)B_X.$$

Then, the covering number

$$N(K, D) \leq \frac{(1 + \frac{1}{2}\epsilon)^n}{(\frac{1}{2}\epsilon)^n}$$

Corollary 6. Let $\epsilon \in (0, 1)$. For n -dimensional Banach spaces, the unit ball has an ϵ -net of cardinality $(\frac{3}{\epsilon})^n$.

Remark: We can follow the same procedure to formulate a statement for the unit sphere.

Proposition 7 (Discretization of $\|A\|$). Let \mathcal{N} be a $\frac{1}{2}$ -net of S^{n-1} . Then $\|A\| \leq 2 \max_{x \in \mathcal{N}} \|Ax\|_2$. (For simplicity, here we let $\epsilon = \frac{1}{2}$)

Proof. Note that every $z \in S^{n-1}$ can be written as $z = x + u$, where $x \in \mathcal{N}$, $\|u\| \leq \epsilon = \frac{1}{2}$. Then, by the triangle inequality,

$$\|A\| \leq \max_{x \in \mathcal{N}} \|Ax\|_2 + \max_{u: \|u\|_2 \leq \frac{1}{2}} \|Au\|_2$$

But $\max_{u: \|u\|_2 \leq \frac{1}{2}} \|Au\|_2 = \frac{1}{2}\|A\|$ by definition. So we have

$$\|A\| \leq \max_{x \in \mathcal{N}} \|Ax\|_2 + \frac{1}{2}\|A\|,$$

and therefore,

$$\|A\| \leq 2 \max_{x \in \mathcal{N}} \|Ax\|_2$$

□

Note: (Exercise) Further, we can discretize the Euclidean norm

$$\|Ax\|_2 = \max_{y \in S^{m-1}} \langle Ax, y \rangle \leq 2 \max_{y \in \mathcal{M}} \langle Ax, y \rangle,$$

where \mathcal{M} is the $\frac{1}{2}$ -net of S^{m-1} .

Corollary 8. Let \mathcal{N}, \mathcal{M} be $\frac{1}{2}$ -nets of S^{n-1}, S^{m-1} , respectively. Then,

$$\|A\| \leq 4 \max_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \langle Ax, y \rangle.$$

Note: We can find that $|\mathcal{N}| \leq 6^n$ and $|\mathcal{M}| \leq 6^m$ by Corollary 6 with $\epsilon = \frac{1}{2}$.

3 Proof of the Main Theorem

Now, we are ready to prove the main theorem that gives an upper estimate of $\|A\|$.

Proof. Let $t > 0$. By Corollary 8, we have

$$\mathbb{P}(\|A\| > t) \leq \mathbb{P}\left(\max_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \langle Ax, y \rangle > \frac{t}{4}\right).$$

Note that if $\max_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \langle Ax, y \rangle > \frac{t}{4}$, then $\exists x \in \mathcal{N}$ and $y \in \mathcal{M}$ such that

$$\langle Ax, y \rangle > \frac{t}{4}.$$

In other words, the event $(\max_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \langle Ax, y \rangle > \frac{t}{4}) \subseteq \bigcup_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} (\langle Ax, y \rangle > \frac{t}{4})$. Hence,

$$\mathbb{P}(\max_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \langle Ax, y \rangle > \frac{t}{4}) \leq \sum_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \mathbb{P}(\langle Ax, y \rangle > \frac{t}{4})$$

Now, fix any $x \in \mathcal{N}$, $y \in \mathcal{M}$, and note that $\langle Ax, y \rangle$ is a random variable.

Claim: $\langle Ax, y \rangle$ is a subgaussian random variable.

Let $x = (x_1, \dots, x_n) \in \mathbb{R}^n$, and let X_{ij} denote the ij^{th} entry in the random matrix A . Then, the coordinates of Ax are

$$\sum_{j=1}^n X_{ij} x_j.$$

Notice that $x \in S^{n-1}$ so $\sum_{j=1}^n x_j = 1$. Also, X_{ij} are independent mean zero subgaussian random variables.

Then, by Theorem 5 from last lecture, the coordinates of Ax are independent subgaussian random variables with mean zero. Then,

$$\langle Ax, y \rangle = \sum_{i=1}^n (Ax)_i y_i,$$

where $\sum_{i=1}^n y_i = 1$. Using the same theorem again, we see that $\langle Ax, y \rangle$ is subgaussian.

Now, knowing that $\langle Ax, y \rangle$ is subgaussian, we have, by definition,

$$\mathbb{P}(\langle Ax, y \rangle > \frac{t}{4}) \leq C \exp(-ct^2).$$

Then,

$$\sum_{\substack{x \in \mathcal{N} \\ y \in \mathcal{M}}} \mathbb{P}(\langle Ax, y \rangle > \frac{t}{4}) \leq |\mathcal{N}| \cdot |\mathcal{M}| \cdot C \exp(-ct^2) \leq 6^{m+n} \cdot C \exp(-ct^2)$$

Thus, with $t = C(\sqrt{n} + \sqrt{m})$, we obtain the desired result. □

Remark:

1. This proof is due to Litvak-Pajar-Rudelson-Tomczak-Vershynin [6].
2. The entries of A must essentially be subgaussian for this result to hold. (Slight fluctuation is allowed)
3. Open problem: (fluctuations) $\mathbb{P}(\|A\| - \mathbb{E}\|A\| > t) \leq ?$
 - Best result by Meckes [4]
 - If A is gaussian, then $\mathbb{E}\|A\| \leq \sqrt{n} + \sqrt{m}$.

References

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