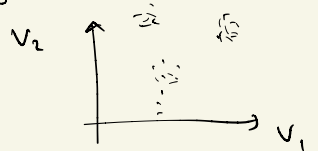


## LECTURE 19

- Last class: How to visualize a high-dim. distribution? (e.g. human genome)  
 $X \in \mathbb{R}^d$  r. vector,  $\mathbb{E}X = 0$ ,  $\text{Cov}(X) = \mathbb{E}XX^T =: \Sigma$ .  
e.g.  $X =$  genome of a random person in the world

Eigenvectors  $v_i$  of  $\Sigma =$  "principal components" of  $X$

- PCA: reduce dimension  $\mathbb{R}^d \rightarrow \mathbb{R}^2$  by projecting  $X$  onto  $\text{span}\{v_1, v_2\}$



- PCA assumes that we can compute the population cov. matrix

$$\Sigma = \text{Cov}(X) = \mathbb{E}XX^T$$

average over  $\uparrow$  population

But we don't have data of all population. We have:

- Finite sample  $X_1, \dots, X_n \in \mathbb{R}^d$  iid copies of  $X$ .  $\Rightarrow$  We approximate  $\Sigma$  by

$$\Sigma_n = \frac{1}{n} \sum_{i=1}^n X_i X_i^T$$

"Sample covariance matrix"

and hope that  $\text{PCA}(\text{sample}) = \text{PCA}(\text{population})$ , i.e.

$$\lambda_i(\Sigma_n) \approx \lambda_i(\Sigma) \quad \text{and} \quad v_i(\Sigma_n) \approx v_i(\Sigma). \quad (*)$$

- How large is  $n$ ?  $n = O(\log d)$ ?  $n = O(d)$ ?  $n = O(e^d)$ ?

### COVARIANCE ESTIMATION PROBLEM

$\uparrow$   
Curse of h.D.?

- OUR GOAL  $n = O(d)$  suffices for  $(*)$ ; smaller for  $\sim$  low rank.

- PLAN
  1. Approximate  $\Sigma_n \approx \Sigma$  in operator norm
  2. Use perturbation theory to conclude  $(*)$

- By **HW** (operator norm for symmetric matrices),

$$\|\Sigma_n - \Sigma\| = \max_{v \in S^{d-1}} \left| \underbrace{v^T (\Sigma_n - \Sigma) v}_{= v^T \Sigma_n v - v^T \Sigma v} \right|$$

where  $S^{d-1}$  = unit sphere in  $\mathbb{R}^d$

$$\bullet \quad v^T \Sigma v = v^T \mathbb{E} X X^T v = \mathbb{E} \underbrace{v^T X}_{\langle X, v \rangle} \underbrace{X^T v}_{\langle X, v \rangle} = \mathbb{E} \langle X, v \rangle^2$$

$$\bullet \quad v^T \Sigma_n v = \frac{1}{n} \sum_{i=1}^n \langle X_i, v \rangle^2$$

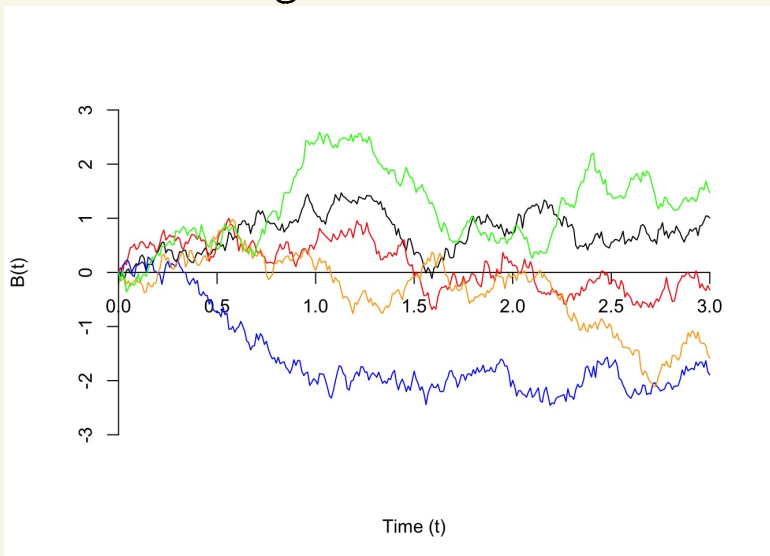
$$\Rightarrow \|\Sigma_n - \Sigma\| = \max_{v \in S^{d-1}} \left| \frac{1}{n} \sum_{i=1}^n \langle X_i, v \rangle^2 - \mathbb{E} \langle X, v \rangle^2 \right|$$

!!  $Z(v)$  random variable.

- $(Z(v))_{v \in S^{d-1}}$  Random process indexed by  $v \in S^{d-1}$ .

- Compare to the Brownian motion (a.k.a. Wiener process)

$(B(t))_{t \in [0, \infty)}$  indexed by time.



$$\bullet \quad \mathbb{E} \max_{t \leq T} |B(t)| \asymp \sqrt{T} \quad \mathbb{E} \max_{v \in S^{d-1}} |Z(v)| \leq ?$$

- Difficulty: a continuum of points in  $S^{d-1}$ ,  $\Rightarrow$  Discretize:

# THE $\epsilon$ -NET METHOD

## Prop (Finding a net)

$\forall \epsilon > 0$ , the unit sphere  $S^{d-1}$  has an  $\epsilon$ -net  $x_1, \dots, x_N$  with

$$\left(\frac{1}{\epsilon}\right)^d \leq N \leq \left(\frac{2}{\epsilon} + 1\right)^d$$

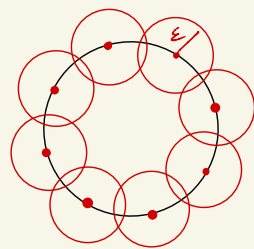
↑ lecture 3

$\Leftrightarrow$  i.e.  $\forall pt \in S^{d-1}$  is within dist.  $\epsilon$  from some  $x_i$ :  
 $\forall x \in S^{d-1} \exists i : \|x - x_i\|_2 \leq \epsilon$

$\Leftrightarrow$   $\epsilon$ -balls centered at  $x_i$  cover  $S^{d-1}$

### Algorithm

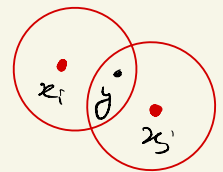
- Choose  $\forall x_1$
- Choose  $\forall x_2$  at dist  $> \epsilon$  from  $x_1$
- Choose  $\forall x_3$  at dist  $> \epsilon$  from  $\{x_1, x_2\}$
- Choose  $\forall x_4$  at dist  $> \epsilon$  from  $\{x_1, x_2, x_3\}$
- ...
- STOP whenever impossible



$\Rightarrow$  COVER

• Claim: The  $(\epsilon/2)$ -balls centered at  $x_i$  are disjoint

If not,  $\exists i \neq j, \exists y : \begin{cases} \|x_i - y\|_2 \leq \epsilon/2 \\ \|x_j - y\|_2 \leq \epsilon/2 \end{cases}$



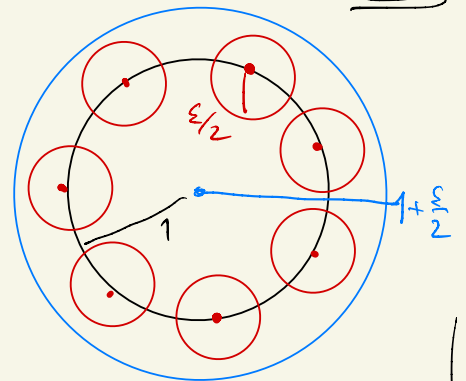
$$\Delta \Rightarrow \|x_i - x_j\|_2 \leq \frac{\epsilon}{2} + \frac{\epsilon}{2} = \epsilon$$

But all  $x_i$  are  $\epsilon$ -separated by construction  $\Leftarrow$

• All these balls lie in the  $(1 + \epsilon/2)$ -ball centered at 0  $\Rightarrow$

$$\text{Vol}(B(1 + \epsilon/2)) \geq N \cdot \text{Vol}(B(\epsilon/2))$$

$$\Rightarrow N \leq \frac{\text{Vol}(B(1 + \epsilon/2))}{\text{Vol}(B(\epsilon/2))} = \left(\frac{1 + \epsilon/2}{\epsilon/2}\right)^d = \left(\frac{2}{\epsilon} + 1\right)^d$$



Remark Covering  $\approx$  packing.

Prop (Computing the operator norm on a net)

Let  $A$  be an  $m \times n$  matrix,  $\mathcal{N} \subset S^{n-1}$  an  $\varepsilon$ -net. Then

$$\|A\| \leq \frac{1}{1-\varepsilon} \max_{x \in \mathcal{N}} \|Ax\|$$

By def of operator norm,  $\exists u \in S^{n-1}$ :

$$\|Au\|_2 = \|A\|. \quad (*)$$

By def of  $\varepsilon$ -net,  $\exists x \in \mathcal{N}$ :

$$\|x-u\|_2 \leq \varepsilon.$$

$$\begin{aligned} \Rightarrow \|Ax - Au\|_2 &= \|A(x-u)\|_2 \leq \|A\| \cdot \|x-u\|_2 && (\text{def of operator norm}) \\ &\leq \|A\| \cdot \varepsilon. && (**) \end{aligned}$$

$$\begin{aligned} \Rightarrow \|Ax\|_2 &= \|Au - (Ax - Au)\|_2 \geq \|Au\|_2 - \|Ax - Au\|_2 && (\Delta \text{ ineq.}) \\ &\geq \|A\| - \|A\| \cdot \varepsilon && (\text{by } (*) \text{ and } (**)) \\ &= (1-\varepsilon) \|A\|. \end{aligned}$$