
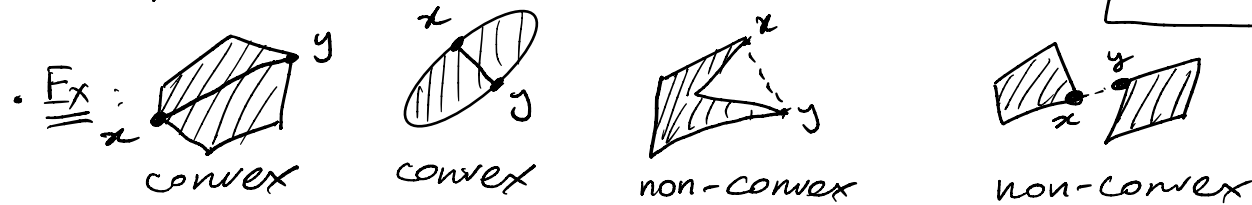


## LECTURE 2

- One more example of how probability can help in high dim's: computational geometry

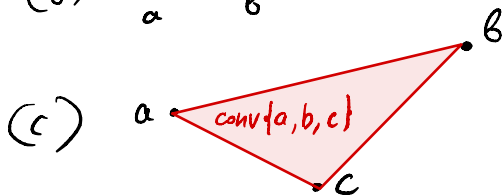
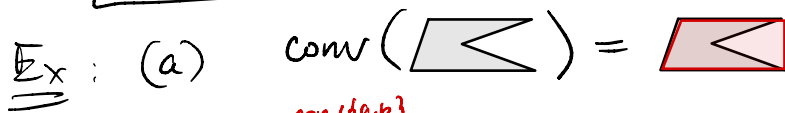
Def A set  $T \subset \mathbb{R}^n$  is convex if  $\forall x, y \in T, [x, y] \in T$

↑  
interval

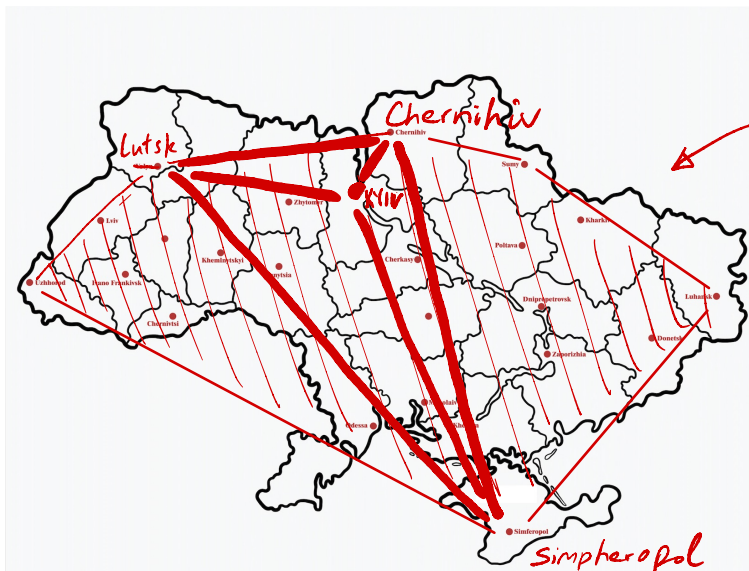



- How can one transform nonconvex set  $\rightarrow$  convex?

Def The convex hull of a set  $T \subset \mathbb{R}^n$ , denoted  $\text{conv}(T)$ , is the smallest convex set that contains  $T$ .



(d)



$T =$  cities in Ukraine  
 $\text{conv}(T)$

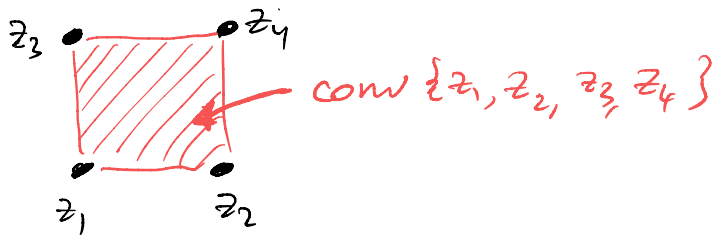
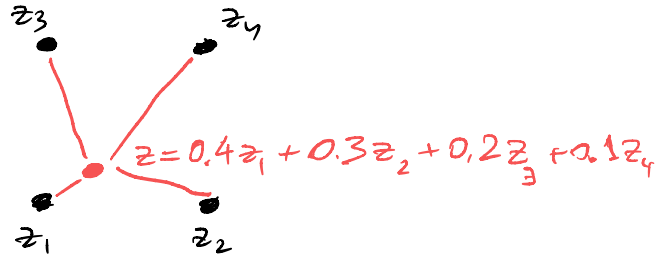
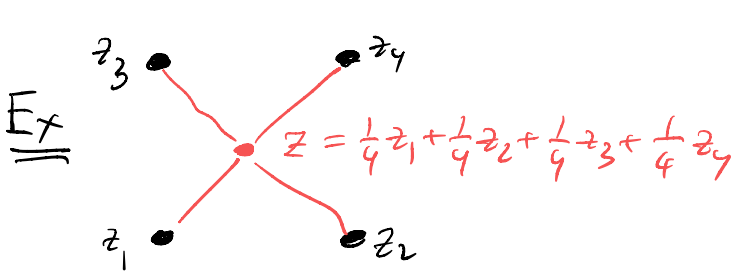
## FACT (CONVEX COMBINATION)

$\forall z \in \text{conv}(T)$  can be expressed as

$$z = \sum_{i=1}^m \lambda_i z_i, \text{ where } \lambda_i \geq 0, \sum_{i=1}^m \lambda_i = 1, z_i \in T \quad (*)$$

• PROOF: DIY (do it yourself).

• (\*) is called a "convex combination". It is a linear combination, similar to a basis expansion of  $z$ , but non-unique (if  $m > n$ )



Q: Computing convex hull: how large is  $m$ ?

Caratheodory Thm  $\forall$  point in  $\text{conv}(T)$  can be expressed as a convex combination of  $\leq n+1$  points from  $T$ .

Ex Kyiv  $\in \text{conv}\{\text{Lutsk, Chernihiv, Simpheropol}\}$

$n=2$ .

Remark,  $n \neq 1$  is unimprovable (e.g. in the Kyiv example above)

But: if we allow to approximate  $x$ , then a dramatic improvement!

$$\max\{\|x-y\|_2 : x, y \in T\}$$

Thm (Approx. Carathéodory) Let  $T \subset \mathbb{R}^n$ ,  $\text{diam}(T) \leq 1$ .

Then  $\forall x \in \text{conv}(T)$ ,  $\forall k \in \mathbb{N} \quad \exists x_1, \dots, x_k \in T$ :

$$\left\| x - \frac{1}{k} \sum_{i=1}^k x_i \right\|_2 \leq \frac{1}{\sqrt{2k}}$$

Euclidean norm in  $\mathbb{R}^n$ .

Remark || Dimension-free!

Does NOT depend on  $n$  or geometry of  $T$ .

PROOF by a probabilistic method: the "empirical method of Maurey"

WILL USE STANDARD FACTS OF PROBABILITY:

① Def of expectation of a discrete r.v.  $X$ :

if  $X$  takes values  $x_i$  with prob.  $p_i$ ,

$$\mathbb{E}X \stackrel{\text{def}}{=} \sum_i p_i x_i \quad \left( \text{analogous to continuous case} \right)$$

(where  $\mathbb{E}X = \int_{-\infty}^{\infty} x p(x) dx$ )

②  $\text{Var}(X) = \mathbb{E}(X - \mathbb{E}X)^2 = \mathbb{E}(X^2) - (\mathbb{E}X)^2 \stackrel{(*)}{=} \frac{1}{2} \mathbb{E}(X - X')^2$

where  $X'$  is an independent copy of  $X$

(i.e.  $X, X'$  are independent and have the same distribution)

Proof of (\*):  $\mathbb{E}(X - X')^2 = \mathbb{E}[X^2 - 2XX' + (X')^2]$

$$= \mathbb{E}[X^2] - 2 \underbrace{\mathbb{E}(X)}_{\mathbb{E}X} \underbrace{\mathbb{E}(X')}_{\mathbb{E}X} + \mathbb{E}(X')^2 \quad (\text{linearity \& independence})$$

$$= 2 \left( \mathbb{E}(X^2) - (\mathbb{E}X)^2 \right) = 2 \text{Var}(X). \quad \text{QED}$$

③ Extend (\*) for random vectors  $X$  in  $\mathbb{R}^n$ :

$$\mathbb{E}\|X - \mathbb{E}X\|_2^2 = \frac{1}{2} \mathbb{E}\|X - X'\|_2^2 \quad \text{where } X' \text{ is an indep. copy of } X \text{ (HW2)}$$

# Proof of Approximate C.T.

- Fix  $x \in \text{conv}(T)$ , express it as a convex combination

$$x = \sum_{i=1}^m \lambda_i z_i, \quad \lambda_i \geq 0, \quad \sum_{i=1}^m \lambda_i = 1, \quad z_i \in T.$$

*potentially large*

- Interpret  $\lambda_i$  as probabilities, the sum as expectation

let r. vector  $Z$  take value  $z_i$  with prob.  $\lambda_i$

$$\Rightarrow x = \mathbb{E}Z$$

- Consider indep. copies  $Z_1, Z_2, \dots$  of  $Z$ .

$$\text{SLLN: } \frac{1}{k} \sum_{i=1}^k Z_i \rightarrow \mathbb{E}Z = x \quad \text{a.s.} \quad P(\mathbb{E})$$

- Error:

$$\mathbb{E} \left\| x - \frac{1}{k} \sum_{i=1}^k Z_i \right\|_2^2 = \mathbb{E} \left\| \frac{1}{k} \sum_{i=1}^k (Z_i - \underbrace{x}_{\mathbb{E}Z_i}) \right\|_2^2$$

$$x = (x_1, \dots, x_n) \in \mathbb{R}^n$$

$$\|x\|_2^2 = x_1^2 + \dots + x_n^2$$

$$= \frac{1}{k^2} \sum_{i=1}^k \underbrace{\mathbb{E} \|Z_i - \mathbb{E}Z_i\|_2^2}_{\mathbb{E} \|Z - \mathbb{E}Z\|_2^2} \quad (\text{variance of sum} = \text{sum of variances, for vectors})$$

HW 2

$$= \frac{1}{k} \mathbb{E} \|Z - \mathbb{E}Z\|_2^2 = \frac{1}{2k} \mathbb{E} \|Z - Z'\|_2^2 \quad (\text{fact } \textcircled{3} \text{ previous page})$$

$$\leq \frac{1}{2k} \quad \wedge \quad \text{diam}(T) \leq 1 \quad \text{since } Z, Z' \in T$$

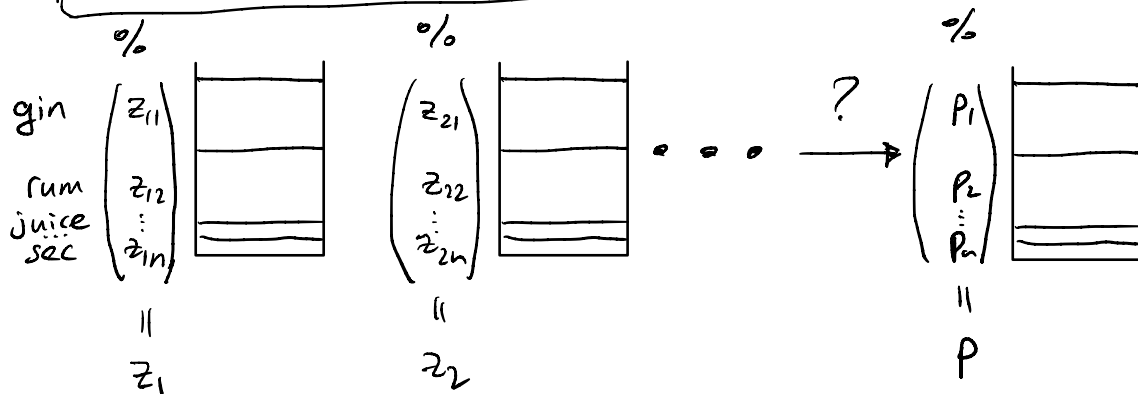
- $\Rightarrow \exists$  realization of the r.v.'s  $Z_1, \dots, Z_k$  s.t.

$$\left\| x - \frac{1}{k} \sum_{i=1}^k Z_i \right\|_2^2 \leq \frac{1}{2k} \quad \text{Since } Z_i \in T, \quad \text{QED}$$



## Application of Approx. Caratheodory Thm:

Cocktail Problem You are given  $N$  glasses with different cocktails, each made by mixing  $n$  ingredients in different proportions. Make a glass of cocktail with given proportions  $p_1, \dots, p_n$ .



Equivalently, we need to find  $\lambda_1$  (portion of glass 1),  $\lambda_2$  (port. of glass 2), ...

such that:  $p = \sum_{i=1}^N \lambda_i z_i$ ,  $\lambda_i \geq 0$ ,  $\sum_{i=1}^N \lambda_i = 1$

$\uparrow$  proportions need to be nonnegative  
 $\uparrow$  and sum to 100%

- i.e. we need to express  $p$  as a convex combination of vectors  $z_1, \dots, z_N$
- "Convex program" finds a solution in polynomial time.
- Approx. Caratheodory Thm transforms it into an approximate solution with few glasses, mixed in equal proportions;
  - $\uparrow$  independent of  $n, N$
  - a fast randomized alg.

# Relevance of cocktail problem

## (a) (Portfolio building)

ingredients = stocks

glasses of cocktails = mutual funds

empty glass = portfolio

Problem: create a new mutual fund with a given combination of stocks by combining the mutual funds that are available on the market.

Solution: as above - a fast randomized algorithm builds a portfolio with few mutual funds.

## (b) (Factor analysis)

$z_1, \dots, z_N$  = a dictionary of factors

that need to explain  $z = \underline{\text{behavior}}$  (consumer, animal, etc.)

Sol: behavior is explained by few factors.

$z = \sum \lambda_i z_i \Rightarrow$  factor 1 explains  $\lambda_1\%$  of behavior,  
...

$\Rightarrow$  A PARSIMONIOUS MODEL.