

LECTURE 37

Applications of Sauer-Shelah Lemma

- ① Consider any n -point set $X \subset \mathbb{R}^2$.
How many subsets of X are cut by circles?

Prop At most $\sum_{k=0}^3 \binom{n}{k} \approx \frac{n^3}{6}$ for large n

Proof: Consider the class of functions

$$\mathcal{F} = \{ \mathbb{1}_C : C \text{ is a circle in } \mathbb{R}^2 \}$$

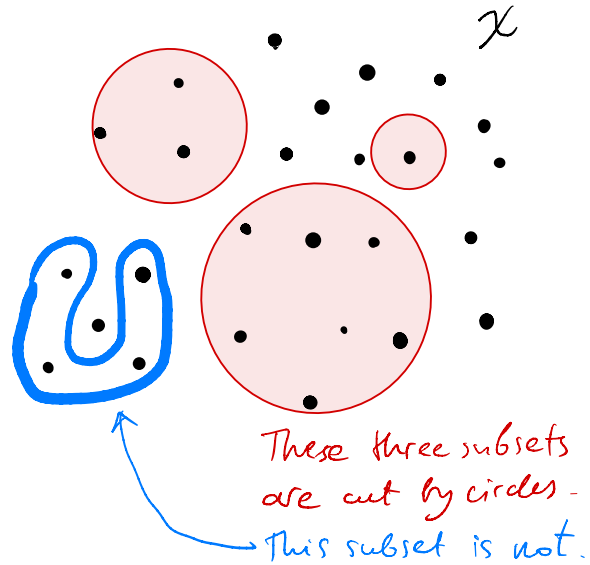
$$vc(\mathcal{F}) = 3.$$

Restrict the domain of each function $f \in \mathcal{F}$ to X

$vc(\mathcal{F}|_X) \leq vc(\mathcal{F}) = 3$. Sauer-Shelah Lemma \Rightarrow

$$|\mathcal{F}|_X| \leq \sum_{k=0}^3 \binom{n}{k}.$$

Remark Prop is optimal for most sets X
(whenever no 4 pts lie on the same circle)

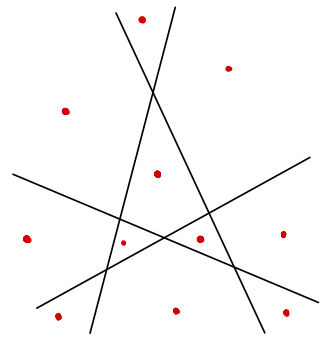


$\{ \mathbb{1}_{C \cap X} : C \text{ is a circle} \}$
indicators of the
subsets of X cut by
circles

② Consider $\forall n$ lines in \mathbb{R}^2 .

How many regions do they partition \mathbb{R}^2 into?

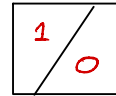
$$\underline{\text{Prop}} \quad \text{At most } \sum_{k=0}^2 \binom{n}{k} = 1 + \frac{n(n+1)}{2}.$$



These four lines break \mathbb{R}^2 into 11 regions

Ex: $n=4$ lines $\Rightarrow 1 + \frac{4 \cdot 5}{2} = 11$ regions

Proof • Choose the orientation of each line \rightarrow



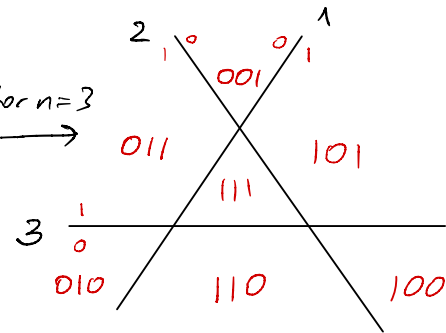
\Rightarrow regions are encoded as binary strings of length n

e.g. for $n=3$

• Identify a binary string with a Boolean function on $\{1, \dots, n\}$

e.g. $0101 \rightarrow f(1)=0, f(2)=1, f(3)=0, f(4)=1$.

Call the class of these functions \mathcal{F} .



• Then $|\text{regions}| = |\mathcal{F}| \leq \sum_{k=0}^d \binom{n}{k}$ where $d = \text{vc}(\mathcal{F})$

Sauer-Shelah

• Claim: $\text{vc}(\mathcal{F}) \leq 2$. Assume the contrary: $\text{vc}(\mathcal{F}) \geq 3$

$\Rightarrow \exists$ a 3-element subset of $\{1, \dots, n\}$ shattered by \mathcal{F}

$\Leftrightarrow \exists$ 3 lines that realize all 8 binary strings of length 3

\Rightarrow these 3 lines break \mathbb{R}^2 into at least 8 regions.

But this is impossible: $\max \# \text{regions} = 7$. QED.

⑥ \exists a simple inductive proof of Prop.

Remarks ① More generally, $\forall n \geq d$ hyperplanes partition \mathbb{R}^d into at most $\sum_{k=0}^d \binom{n}{k}$ regions [Buck 1943]. The proof is similar.

② The bound is optimal, attained \forall hyperplanes in general position.

③ If not in general position, see enumerative combinatorics / hyperplane arrangements / Möbius function / Zaslavsky formula.

EMPIRICAL PROCESSES

- Let X be a r.v. taking values in \mathcal{X} set \mathcal{X} ,
 X_1, \dots, X_n be independent copies of X .
- Law of large numbers $\Rightarrow \forall$ Boolean function $f: \mathcal{X} \rightarrow \{0,1\}$:

$$\frac{1}{n} \sum_{i=1}^n f(X_i) \xrightarrow{\text{a.s.}} \mathbb{E} f(X) \quad \text{as } n \rightarrow \infty$$

• Deviation:

$$\begin{aligned} \mathbb{E} \left| \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E} f(X) \right| &\stackrel{L^1 \leq L^2}{\leq} \left(\mathbb{E} \left| \dots \right|^2 \right)^{1/2} = \text{Var} \left(\frac{1}{n} \sum_{i=1}^n f(X_i) \right)^{1/2} \\ &= \left[\frac{1}{n^2} \sum_{i=1}^n \underbrace{\text{Var}(f(X_i))}_{\substack{\wedge \\ 1 \text{ (f Boolean)}}} \right]^{1/2} = \frac{1}{\sqrt{n}} \quad (\text{"Weak LLN"}) \end{aligned}$$

- Is this true uniformly over all Boolean functions f ?

$$\mathbb{E} \sup_{f \in \mathcal{F}} \left| \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E} f(X) \right| \rightarrow 0 \quad \text{as } n \rightarrow \infty ?$$

NO For $f = 1_{\{X_1, \dots, X_n\}}$, $\frac{1}{n} \sum_{i=1}^n f(X_i) = 1$ but $\mathbb{E} f(X) = 0$



- But it is true uniformly over $f \in \mathcal{F}$ whenever $vc(\mathcal{F}) < \infty$:

THM (Uniform Law of Large Numbers)

If \mathcal{F} is a Boolean class with $d = vc(\mathcal{F}) < \infty$, then

$$\mathbb{E} \sup_{f \in \mathcal{F}} \left| \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E} f(X) \right| \leq C \sqrt{\frac{d \log n}{n}}.$$

Remarks ① Same rate $O(1/\sqrt{n})$ as in WLLN!

② $\left(\frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E} f(X) \right)_{f \in \mathcal{F}}$ is called an *empirical process*.

The proof uses a new tool: — 3 —