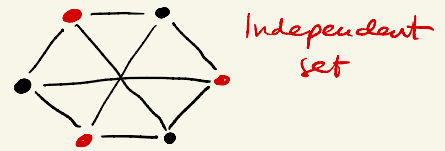


### ③ Turan's Theorem '1941 [AS]

"If in a group of 100 people each person knows 3 other people, then there are 25 people who don't know each other".

More generally:

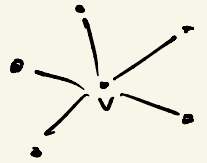
Thm  $\forall$   $d$ -regular graph on  $n$  vertices  $G(V, E)$   
 $\Rightarrow$  an independent set of cardinality  $\geq \frac{n}{d+1}$



every vertex has exactly  $d$  neighbors  
 a subset of vertices no two of which are neighbors

Proof Choose a random ordering of  $V$  (uniformly)

- $I := \{v \in V : v \text{ is the least element among } v \text{ and its neighbors}\}$



$$\mathbb{E}|I| = \mathbb{E} \sum_{v \in V} \mathbb{1}_{\{v \in I\}} = \sum_{v \in V} \underbrace{P\{v \in I\}}_{\frac{1}{d+1}} = \frac{n}{d+1}$$

- $I$  is an independent set:  $\frac{1}{d+1}$  (prob. that  $v$  is the least among  $d+1$  elements in random order)  
 if  $u, v \in I$  and  $(u, v) \in E$  then  $u < v$  and  $v < u \Rightarrow \nabla \quad \square$

Remark: Turan's Bound is optimal for the union of  $\frac{n}{d+1}$  cliques of size  $d+1$ .

④ k-SAT problem

- A Boolean formula on  $n$  literals is a function  $\Phi: \{T, F\}^n \rightarrow \{T, F\}$
- $\Phi$  is satisfiable if  $\exists$  an assignment  $x$  such that  $\Phi(x) = T$
- $\forall$  Boolean formula can be expressed in the conjunctive normal form (CNF):  
AND of ORs (of  $x_i$ 's and  $\bar{x}_i$ 's), e.g.

$$\Phi(x) = (x_1 \vee \bar{x}_2) \wedge (x_1 \vee \bar{x}_2 \vee x_3) \wedge \bar{x}_1.$$

OR
AND
negation

Even 3-sat  $\in$  NP

Q Consider a random "k-SAT" formula with  $rn$  clauses each of length  $k$ , where each literal is chosen independently uniformly from  $\{x_1, \dots, x_n, \bar{x}_1, \dots, \bar{x}_n\}$ . What is the probability that  $\Phi$  is satisfiable?

$$\Phi(x) = (\underbrace{? \vee ? \vee \dots \vee ?}_k) \wedge (\underbrace{? \vee ? \vee \dots \vee ?}_k) \wedge \dots \wedge (\underbrace{? \vee ? \vee \dots \vee ?}_k)$$

$\swarrow$   $rn$  clauses  $\searrow$   
 $\uparrow$  each "?" can be a literal  $x_i$  or its negation  $\bar{x}_i$

Thm If  $r > 2^k \ln 2$  and  $n \rightarrow \infty$  (while  $r, k$  are fixed), then  $P\{\Phi \text{ is satisfiable}\} \rightarrow 0$

Proof  $P\{\Phi \text{ is satisfiable}\} = P\{\exists x \in \{T, F\}^n : \Phi(x) = T\}$   
 $\leq \sum_{x \in \{T, F\}^n} P\{\Phi(x) = T\} \quad (\text{union bound})$

$$\Phi(x) = \Phi_1(x) \wedge \dots \wedge \Phi_{rn}(x) \Rightarrow$$

$$P\{\Phi(x) = T\} = P\{\Phi_i(x) = 1 \forall i=1, \dots, rn\} \stackrel{\text{indep}}{=} \prod_{i=1}^{rn} P\{\Phi_i(x) = 1\} = \left(1 - \frac{1}{2^k}\right)^{rn}$$

$$\stackrel{\ominus}{=} 2^n \left(1 - \frac{1}{2^k}\right)^{rn} \leq \underbrace{\left(2e^{-r/2^k}\right)^n}_{\uparrow \text{ if } r > 2^k \ln 2} \rightarrow 0 \text{ as } n \rightarrow \infty. \quad \square$$

$\parallel$   
 $\uparrow - \left(\frac{1}{2}\right)^k$  (since  $\Phi_k = k$  or's)

Remark If  $k \geq 3$ ,  $r < 2^k \ln 2 - k$ ,  $P\{\Phi \text{ is satisfiable}\} \rightarrow 1$  ( $n \rightarrow \infty$ ) [Alon-Naor-Peres 05]

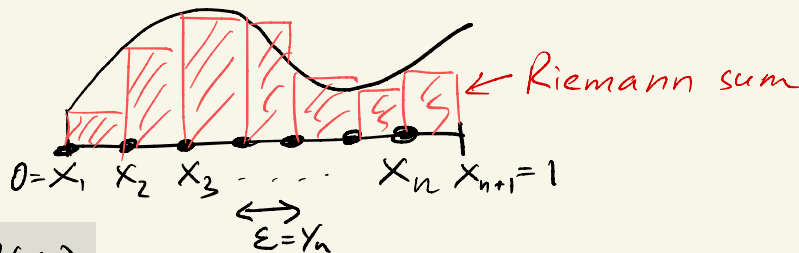
# ⑤ Monte-Carlo Integration

Problem For a given function  $f: [0,1]^d \rightarrow \mathbb{R}$ , numerically compute

$$\int_0^1 \dots \int_0^1 f(x_1, \dots, x_d) dx_1 \dots dx_d = \int_{[0,1]^d} f(x) dx$$

## METHOD 1

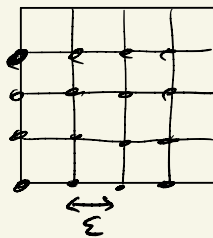
• If  $d=1$ : use the grid



$$\int_0^1 f(x) dx \approx \sum_{i=1}^n \underbrace{(x_{i+1} - x_i)}_{\epsilon} f(x_i) = \frac{1}{n} \sum_{i=1}^n f(x_i)$$

"Resolution"  $\epsilon = 1/n$

• If  $d=2$ , do similarly but use the grid



$$\Rightarrow n = \left(\frac{1}{\epsilon}\right)^2 \text{ pts.}$$

• For general dimension  $d$ , we need  $n = \left(\frac{1}{\epsilon}\right)^d \text{ pts.}$

Exponential in  $d$ . Too large.

"THE CURSE OF DIMENSIONALITY"



Method 2: Monte-Carlo

- Instead of choosing  $x_i$  on the grid, choose them at random independently:  $X_i \sim \text{Unif}([0,1]^d)$

$\Rightarrow f(x_i)$  are i.i.d. r.v.'s.

$$\frac{1}{n} \sum_{i=1}^n f(x_i) \stackrel{?}{\approx} \int_{[0,1]^d} f(x) dx$$

- $$\mathbb{E} \left[ \frac{1}{n} \sum_{i=1}^n f(x_i) \right] = \frac{1}{n} \sum_{i=1}^n \underbrace{\mathbb{E} f(x_i)}_{\mathbb{E} f(x) \text{ by identical distribution}} = \mathbb{E} f(x)$$

$$= \int_{\mathbb{R}^d} f(x) p(x) dx, \text{ where density is } p(x) = \begin{cases} 1, & x \in [0,1]^d \\ 0, & \text{---} \end{cases}$$

$$= \int_{[0,1]^d} f(x) dx. \quad \text{😊}$$

$\Rightarrow$  we have an unbiased estimator of the integral

- Rate of convergence?  $L^2$  error ("MSE"):

$$\mathbb{E} \left( \underbrace{\frac{1}{n} \sum_{i=1}^n f(x_i)}_{\bar{z}} - \underbrace{\int_{[0,1]^d} f(x) dx}_{\mathbb{E} z} \right)^2 = \text{Var} \left( \frac{1}{n} \sum_{i=1}^n f(x_i) \right)$$

$$= \frac{1}{n^2} \sum_{i=1}^n \underbrace{\text{Var}(f(x_i))}_{\text{Var}(f(x)) \text{ by identical distribution}} = \frac{\text{Var}(f(x))}{n}$$

$$\leq \frac{1}{n} \quad \text{e.g. if } |f(x)| \leq 1 \quad \forall x.$$

- Chebyshev's inequality  $\Rightarrow$  with probability  $\geq 0.99$ , 
$$\left| \frac{1}{n} \sum_{i=1}^n f(x_i) - \int_{[0,1]^d} f(x) dx \right| \leq \frac{10}{\sqrt{n}} \quad \text{😊}$$

Regardless of dimension  $d$ !