

Example (maximum of Gaussians) Let g_1, g_2, \dots be iid $N(0,1)$ r.v.'s. Then

$$\limsup_n \frac{g_n}{\sqrt{2 \ln n}} = 1 \quad \text{a.s.}$$

Proof The conclusion is equivalent to the following statement. $\forall \epsilon > 0$:

$$P \{ g_n > \sqrt{(1+\epsilon) 2 \ln n} \text{ i.o.} \} = 0 \quad (1)$$

AND $P \{ g_n > \sqrt{(1-\epsilon) 2 \ln n} \text{ i.o.} \} = 1 \quad (2)$

To prove this, use Gaussian tail bounds (p.28): if $g \sim N(0,1)$, then

$$\left(\frac{1}{t} - \frac{1}{t^3} \right) \frac{1}{\sqrt{2\pi}} e^{-t^2/2} \leq P\{g > t\} \leq \frac{1}{t\sqrt{2\pi}} e^{-t^2/2} \quad \forall t > 0$$

In particular, $\forall \epsilon > 0 \exists t(\epsilon)$ s.t.:

$$e^{-(1+\frac{\epsilon}{2})t^2/2} \leq P\{g > t\} \leq e^{-t^2/2} \quad \forall t > t(\epsilon)$$

Hence:

- $P \{ g_n > \sqrt{(1+\epsilon) 2 \ln n} \} \leq \exp(- (1+\epsilon) \ln n) = n^{-(1+\epsilon)}$ is summable

\Rightarrow (1) follows from Borel-Cantelli I.

- $P \{ g_n > \sqrt{(1-\epsilon) 2 \ln n} \} \geq \exp \left(- \underbrace{(1+\frac{\epsilon}{2})(1-\epsilon)}_{\parallel} \ln n \right) \leq n^{-(1-\frac{\epsilon}{2})}$ is NOT summable

\Rightarrow (2) follows from Borel-Cantelli II.

□

Example ("Runs") An infinite monkey is hitting keys at random.

How much of William Shakespeare's work will be typed by time n ?

Formally: fix an infinite word $W = (w_1, w_2, \dots)$

e.g. $W = \text{"TOBEORNOTTOBE..."}$

and a random word $X = (X_1, X_2, \dots)$ with all $X_i \sim \text{Unit}(\text{English alphabet})$ independent

e.g. $X = \text{"ABRATOCADATOBESMTHTOBEORIWANTABANANA..."}$
run of length 2 run of length 4 independent
1 2 3 4 5 6 7 8 n $L_n = 6$
 $l_6 = 2, l_7 = 0$

$l_k := \max \{ r : X_k = W_k, X_{k-1} = W_{k-1}, \dots, X_{k-r+1} = W_{k-r+1} \}$ (length of a run at time k)

$L_n := \max_{1 \leq k \leq n} l_k$ (max length of a run by time n)

Thm $\frac{L_n}{\log_{26} n} \rightarrow 1$ a.s.

(26 = size of English alphabet)

Proof Upper bound:

$P\{l_n \geq r\} = P\{r \text{ last letters are correct}\} = \frac{1}{26^r}$

$\Rightarrow P\{l_n > (1+\epsilon) \log_{26} n\} \leq \frac{1}{26^{(1+\epsilon) \log_{26} n}} = \frac{1}{n^{1+\epsilon}}$ is summable

\Rightarrow by Borel-Cantelli I,

$P\{l_n > (1+\epsilon) \log_{26} n \text{ i.o.}\} = 0$

\Rightarrow with probability 1, *possibly random*

$\exists m \in \mathbb{N} \forall n \geq m: l_n < (1+\epsilon) \log_{26} n.$ } $\Rightarrow l_n < (1+\epsilon) \log_{26} n + m$

Moreover, $\forall n < m: l_n \leq n < m$ trivially

$\Rightarrow L_n \leq (1+\epsilon) \log_{26} n + m$

$\Rightarrow \frac{L_n}{\log_{26} n} \leq 1 + \epsilon + \frac{m}{\log_{26} n} \xrightarrow{0 \text{ as } n \rightarrow \infty} \Rightarrow \limsup_n \frac{L_n}{\log_{26} n} \leq 1 + \epsilon \quad \forall \epsilon > 0$

$\Rightarrow \limsup_n \frac{L_n}{\log_{26} n} \leq 1$ a.s. \odot

Lower bound :

Partition $\{1, \dots, n\}$ into $\frac{n}{(1-\varepsilon)\log_{26} n}$ blocks of length $(1-\varepsilon)\log_{26} n$. Then

$$P\left\{\frac{L_n}{\log_{26} n} < 1-\varepsilon\right\} \leq P\left\{\begin{array}{l} \text{none of the blocks} \\ \text{is a run} \end{array}\right\} \quad \left(\begin{array}{l} \uparrow \\ \text{if some block is a run,} \\ L_n > (1-\varepsilon)\log_{26} n \end{array}\right)$$

$$\leq \left(1 - \frac{1}{26^{(1-\varepsilon)\log_{26} n}}\right)^{\frac{n}{(1-\varepsilon)\log_{26} n}} \quad \leftarrow \# \text{ blocks (independent)}$$

↑
Prob (given block = run)

$$= \left(1 - \frac{1}{n^{1-\varepsilon}}\right)^{\frac{n}{(1-\varepsilon)\log_{26} n}}$$

$$\leq \exp\left(-\frac{n}{n^{1-\varepsilon}(1-\varepsilon)\log_{26} n}\right) \quad (1-x \leq e^{-x})$$

$$= \exp\left(-\frac{n^\varepsilon}{(1-\varepsilon)\log_{26} n}\right) \quad \text{is summable.}$$

\Rightarrow by Borel-Cantelli I,

$$P\left\{\frac{L_n}{\log_{26} n} < 1-\varepsilon \text{ i.o.}\right\} = 0$$

\Rightarrow with probability 1, $\exists m \forall n \geq m: \frac{L_n}{\log_{26} n} \geq 1-\varepsilon$

$$\Rightarrow \liminf \frac{L_n}{\log_{26} n} \geq 1-\varepsilon \Rightarrow \boxed{\liminf \frac{L_n}{\log_{26} n} = 1 \text{ a.s.}}$$

☺

□

Saint Petersburg Paradox II

Consider an ∞ sequence of games in which one

loses $\$2^n$ with prob. $\frac{1}{2^{n+1}}$,

wins $\$1$ with prob. $\frac{2^n}{2^{n+1}}$, $n=1, 2, 3, \dots$

• $X_n :=$ winnings of n 'th game. $\mathbb{E}X_n = -2^n \cdot \frac{1}{2^{n+1}} + 1 \cdot \frac{2^n}{2^{n+1}} = 0$
 $X :=$ total winnings $\Rightarrow \mathbb{E}X = \sum_{n=1}^{\infty} \mathbb{E}X_n = 0$ (*)

• $P(\text{losses occur } \infty \text{ often}) = 0$ by Borel-Cantelli

$\left(E_n := \text{"n-th game is lost"} \Rightarrow \sum_{n=1}^{\infty} \underbrace{P(E_n)}_{\frac{1}{2^{n+1}}} < \infty \right)$

Hence: with probability 1, starting from some game

we will keep winning and will never lose again. \Rightarrow winnings = (∞)

But our expected winnings = (0) !?

Resolution of paradox: $\mathbb{E}X$ does NOT exist;

the limit in (*) is not justified

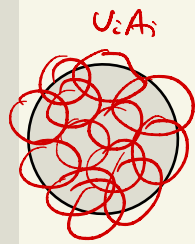
Example (\sim E. Stein's covering lemma)

Thm let $A_1, A_2, \dots \subset S^{n-1}$ be any measurable subsets such that

$$\sum_{i=1}^{\infty} \mu(A_i) = 0.$$

↑
"surface area"

Then $\exists U_1, U_2, \dots \in O(n)$ such that μ -almost every point $x \in S^{n-1}$ belongs to ∞ many sets $U_i A_i$.



Proof let $U_1, U_2, \dots \in \text{Unif}(O(n))$ and $X \sim \text{Unif}(S^{n-1})$, all independent.
↑ Haar measure ↑ μ , prob. meas. wlog

Then $U_i^{-1} X, U_2^{-1} X, \dots \in \text{Unif}(S^{n-1})$ independent (check!) (*)

Consider events

$$E_i := \{U_i^{-1} X \in A_i\} = \{X \in U_i A_i\}$$

Note that E_i are independent & $P(E_i) = \mu(A_i)$ by (*).

By assumption, $\sum_{i=1}^{\infty} P(E_i) = \infty$.

Borel-Cantelli II \Rightarrow

$$1 = P\{E_i \text{ occur i.o.}\} = E \mathbb{1}_{\{X \in U_i A_i \text{ i.o.}\}}$$

$$\Rightarrow \mathbb{1}_{\{X \in U_i A_i \text{ i.o.}\}} = 1 \text{ a.s. (**)}$$

↑
w.r. to U_i 's and X

$\Rightarrow \exists U_i : (**)$ holds a.s. w.r. to X (check!)

□