MATH 13 SAMPLE FINAL EXAM SOLUTIONS

WINTER 2014

Problem 1 (15 points). For each statement below, circle T or F according to whether the statement is true or false. You do NOT need to justify your answers.

- T (F) Every countably infinite subset of \mathbb{R} is well-ordered.
- T (F) For all integers a, b, and c, if $a \mid 2b$ and $b \mid 2c$ then $a \mid 2c$.
- (T) F If A is a set and $\mathcal{P}(A)$ is infinite, then A must also be infinite.
- (I) F There is a bijection from the closed interval $[0, \infty)$ to the open interval $(0, \infty)$.
- \bigcirc F For all sets A and B, if there is a surjection from A to B and A is countable, then B is countable.
- ① F If R is a symmetric relation on the set A, then $(A \times A) R$ must also be a symmetric relation on A.
- T F There are sets A, B, and C and functions $f: A \to B$ and $g: B \to C$ such that f is not injective and $g \circ f: A \to C$ is injective.
- T (F) For every set A, there is a bijection from $\mathcal{P}(A)$ to $A^{\{0,1\}}$.
- \bigcirc F If R is a reflexive relation on the set A, then R^{-1} must also be a reflexive relation on A.
- ① F If R_1 and R_2 are equivalence relations on the set A, then $R_1 \cap R_2$ must also be an equivalence relation on A.
- ① F For all sets A, B, and C, there is a bijection from $A \times (B \times C)$ to $(A \times B) \times C$.
- T $(P \land Q)$ is logically equivalent to $\sim P \land \sim Q$
- $(T) \quad F \quad \emptyset \subseteq \emptyset$
- T For all integers a, b, n with $n \ge 2$, if $ab \equiv 0 \pmod{n}$ then $a = 0 \pmod{n}$ or $b = 0 \pmod{n}$.

Remark. Andrés pointed out an error in my first draft of the solutions. The answer to the last T/F question above is 'F' because the empty set is a well-ordered set of real numbers without a least element—it has no nonempty subsets, so it satisfies the definition of "well-ordered" vacuously. (For the exam I will try to avoid asking questions that hinge on trivial cases.)

Problem 2 (5 points). Let A and B be sets. Prove that the following equality holds:

$$(A \cup B) - (A \cap B) = (A - B) \cup (B - A).$$

Solution.

- (\subseteq) : Let $x \in (A \cup B) (A \cap B)$, so $x \in A \cup B$ and $x \notin A \cap B$.
 - Case 1: Assume $x \in A$. Then $x \notin B$ (or else x would be in $A \cap B$) so $x \in A B$.
 - Case 2: Assume $x \in B$. Then $x \notin A$ (or else x would be in $A \cap B$) so $x \in B A$.

We have shown that, in either case, $x \in (A - B) \cup (B - A)$.

- (\supseteq) : Let $x \in (A B) \cup (B A)$.
 - Case 1: Assume $x \in A B$; that is, $x \in A$ and $x \notin B$. Then $x \in A \cup B$ and $x \notin A \cap B$.
 - Case 2: Assume $x \in B A$; that is, $x \in B$ and $x \notin A$. Then $x \in A \cup B$ and $x \notin A \cap B$.

We have shown that, in either case, $x \in (A \cup B) - (A \cap B)$.

Solution (alternative).

$$(A \cup B) - (A \cap B) = (A \cup B) \cap \overline{(A \cap B)}$$

$$= (A \cup B) \cap (\overline{A} \cup \overline{B}) \quad \text{(De Morgan's law)}$$

$$= ((A \cup B) \cap \overline{A}) \cup ((A \cup B) \cap \overline{B}) \quad \text{(distributive law)}$$

$$= ((A \cap \overline{A}) \cup (B \cap \overline{A})) \cup ((A \cap \overline{B}) \cup (B \cap \overline{B})) \quad \text{(distributive law)}$$

$$= (B \cap \overline{A}) \cup (A \cap \overline{B})$$

$$= (B - A) \cup (A - B)$$

$$= (A - B) \cup (B - A)$$

Problem 3 (5 points). Is this compound statement a tautology? Justify your answer.

$$((P \implies Q) \land (\sim\! P \implies Q)) \implies Q$$

Note: " $\sim P \implies Q$ " means " $(\sim P) \implies Q$," not " $\sim (P \implies Q)$."

Solution. Assume that " $(P \implies Q) \land (\sim P \implies Q)$ " is true. So " $P \implies Q$ " and " $\sim P \implies Q$ " are both true.

- Case 1: Assume P. Then because " $P \implies Q$ " is true, Q is true.
- Case 2: Assume $\sim P$. Then because " $\sim P \implies Q$ " is true, Q is true.

We have shown that in all the cases where the hypothesis is true, the conclusion is also true. So the implication is a tautology.

Remark. We could also solve this problem by making a truth table with four rows for the possible combinations of truth values for P and Q. The type of solution above has the advantage that it is more intuitive (in my opinion.) It also might be faster, especially on T/F questions where you are not asked to write out your argument.

Problem 4 (5 points). Let $a, a', b, b', n \in \mathbb{Z}$ with $n \geq 2$. Prove that if $a \equiv a' \pmod{n}$ and $b \equiv b' \pmod{n}$, then $a + b \equiv a' + b' \pmod{n}$.

Solution. Assume that $a \equiv a' \pmod{n}$ and $b \equiv b' \pmod{n}$. So $n \mid (a'-a)$ and $n \mid (b'-b)$. Therefore $n \mid ((a'-a)+(b'-b))=((a'+b')-(a+b))$, so $a+b \equiv a'+b' \pmod{n}$.

Problem 5 (5 points). Define a sequence (a_1, a_2, a_3, \ldots) by $a_1 = 3$, and for all $n \in \mathbb{N}$,

$$a_{n+1} = \begin{cases} 2a_n & \text{if } n \text{ is odd} \\ 4a_n & \text{if } n \text{ is even.} \end{cases}$$

Prove that $a_n \leq 3^n$ for all $n \in \mathbb{N}$.

Proof. We prove this using strong induction on n.

- Base case: $a_1 = 3 \le 3^1$.
- Induction step: Let $n \in \mathbb{N}$ and assume that $a_i \leq 3^i$ for all $i \leq n$. We want to show that $a_{n+1} \leq 3^{n+1}$. If n is odd, then we have

$$a_{n+1} = 2a_n \le 2 \cdot 3^n \le 3 \cdot 3^n = 3^{n+1},$$

as desired. If n is even, then (using the fact that $n-1 \in \mathbb{N}$ and n-1 is odd) we have $a_{n+1}=4a_n=4\cdot 2a_{n-1}=8a_{n-1}\leq 8\cdot 3^{n-1}\leq 9\cdot 3^{n-1}=3^{n+1},$

as desired.

Therefore $a_n \leq 3^n$ for all $n \in \mathbb{N}$.

Problem 6 (5 points). Let A be a set and let $f: A \to A$ be a function such that f(f(x)) = x for all $x \in A$. Define the relation R on A by

$$R = \{(x, y) \in A : y = x \text{ or } y = f(x)\}.$$

Prove that R is an equivalence relation.

Solution.

- We show that R is reflexive. Let $x \in A$. Then $(x, x) \in R$ because x = x.
- We show that R is symmetric. Let $x, y \in A$ and assume that $(x, y) \in R$. That is, y = x or y = f(x). If y = x, then x = y, so $(y, x) \in R$. If y = f(x), then f(y) = f(f(x)) = x, so $(y, x) \in R$ in this case also.
- We show that R is transitive. Let $x, y, z \in A$ and assume that $(x, y), (y, z) \in R$. If x = y or y = z then clearly $(x, z) \in R$. If not, then y = f(x) and z = f(y), so z = f(f(x)) = x and again $(x, z) \in R$.

Problem 7 (5 points). Prove that there is a bijection from $A \times A$ to $A^{\{0,1\}}$.

Solution. Define a function $G: A \times A \to A^{\{0,1\}}$ by $G((x,y)) = \{(0,x), (1,y)\}^{1}$

- We show that G is injective. Let $(x,y),(x',y') \in A \times A$ and assume that G((x,y)) = G((x',y')). That is, $\{(0,x),(1,y)\} = \{(0,x'),(1,y')\}$. Clearly this means that (0,x) = (0,x') and (1,y) = (1,y') and not the other way around. So x = x' and y = y', which means that (x,y) = (x',y').
- We show that G is surjective. Let $f \in A^{\{0,1\}}$. Then we have $G((f(0), f(1))) = \{(0, f(0)), (1, f(1))\} = f$, so f is in the range of G.

Solution (alternative). Because every bijection has an inverse that is also a bijection, it suffices to prove that there is a bijection from $A^{\{0,1\}}$ to $A \times A$.

Define a function $G: A^{\{0,1\}} \to A \times A$ by G(f) = (f(0), f(1)).

- We show that G is injective. Let $f, f' \in A^{\{0,1\}}$ and assume that G(f) = G(f'). That is, (f(0), f(1)) = (f'(0), f'(1)). So f(0) = f'(0) and f(1) = f'(1). Therefore f = f'.
- We show that G is surjective. Let $(x,y) \in A \times A$. Define the function $f \in A^{\{0,1\}}$ by f(0) = x and f(1) = y. Then G(f) = (x,y).

¹Equivalently, we could define the function G((x,y)) in terms of its values: G((x,y))(0) = x and G((x,y))(1) = y.

²Equivalently, we could write $f = \{(0, x), (1, y)\}.$

Problem 8 (5 points). Prove that there is a partition P of \mathbb{N} such that

- (1) P is infinite, and
- (2) every element of P is infinite.

(Hint: You may use the existence of a bijection $f: \mathbb{N} \to \mathbb{N} \times \mathbb{N}$.)

Proof. Take a bijection $f: \mathbb{N} \to \mathbb{N} \times \mathbb{N}$. Define a set P of subsets of N by

$$P = \{A_i : i \in \mathbb{N}\}$$

where

$$A_i = f^{-1}[\{i\} \times \mathbb{N}] = \{x \in \mathbb{N} : \text{the ordered pair } f(x) \text{ has first coordinate } i\}.$$

We show that P is a partition.

- Every element of P is nonempty, because $\{i\} \times \mathbb{N}$ is nonempty and f is surjective.
- The elements of P are pairwise disjoint because $(\{i\} \times \mathbb{N}) \cap (\{j\} \times \mathbb{N}) = \emptyset$ whenever $i \neq j$, and the inverse images of two disjoint sets (under any function) are disjoint.
- $\mathbb{N} = \bigcup_{i \in \mathbb{N}} A_i$ because for all $x \in \mathbb{N}$ we have $x \in A_i$ where i is the first coordinate of the pair f(x).

We check the remaining two conditions.

- P is infinite because the function $\mathbb{N} \to P$ that sends i to A_i is a bijection. (The fact that it is an injection follows from the argument above, which shows that if $i \neq j$ then A_i and A_j are disjoint nonempty sets, hence $A_i \neq A_j$.)
- Every element A_i of P is infinite because the function $A_i \to \mathbb{N}$ that sends x to the second coordinate of f(x) is easily shown to be a bijection....

Remark. The details of this solution take a bit longer to write out than I intended. The main point is that the set $\mathbb{N} \times \mathbb{N}$ has an infinite partition into the infinite sets $\{1\} \times \mathbb{N}$, $\{2\} \times \mathbb{N}$, etc. and we can then use a bijection of \mathbb{N} with $\mathbb{N} \times \mathbb{N}$ to "transfer" this partition over to a partition of \mathbb{N} . Once you come up with this idea (or some other idea) you should just write as much detail as you have time for and not worry about the rest.

Remark. A more number-theoretic proof idea (without using the hint) would be to define, for all $i \in \mathbb{N}$, the set

$$A_i = \{ x \in \mathbb{N} : 2^{i-1} \mid x \text{ and } 2^i \nmid x \},$$

and to show that the sets $A_1, A_2, ...$ are all distinct infinite sets and that the set $P = \{A_1, A_2, ...\}$ is a partition of \mathbb{N} .